



Figure 1: Building the friend candidate graph for a target user. (a) V_p^C = Users in the friend graph that contain the target user’s phone number in their contact book. (b) V_k^F = friends of V_p^C that have at least k friendship connections with them. (c) The friend candidate graph has the node set = $V_p^C \cup V_k^F$. Edge weights encode candidate interactions, e.g., the number of chat messages exchanged, amount of content created, etc.

	V_p^C	$V_p^C \cup V_0^F$	$V_p^C \cup V_k^F$	V_d^C	$V_d^C \cup V_0^F$	$V_d^C \cup V_k^F$
$ V $	36.8	666.8	396.9	47.1	649.4	398.5
$ E $	37.3	4913.4	2524.9	49.4	4491	2372.6
UES	42.6	829.9	629.6	40.52	704.8	537

Table 1: Avg. # of graph nodes $|V|$, of edges $|E|$, and avg. top5 UES for multiple candidate graph variations. In blue: our final node set.

tricking k users into doing that is even more difficult, as an attacker need not only create a clique of fake friends in the network, but also to connect the fake clique to real users, while creating a pattern.

The choice of k represents a trade-off between privacy and utility. In practice, it makes sense to choose k in a way that depends on the friendship graph sparsity but with an element of randomness; we use $k \leftarrow \text{rand}[\max\{k_{med} - \delta, k_{min}, 2\}, \min\{k_{med} + \delta, k_{max}\}]$, where $\delta \geq 2$, and $k_{min}, k_{max}, k_{med}$ are the min, max and median number of friends per user among the original pool of candidates.

We note that Snapchat uses a user’s contact book only if the user has opted-in to share its contents. Furthermore, we utilize Bloom filters [3] to encode and check contact book relationships, rather than explicitly storing the contact books.

4 EXPERIMENTS

We now experimentally study the impact that using the friend candidate graph instead of the actual friend graph may have on utility. We measure utility in terms of the characteristics of the graph and the user engagement score (UES) calculated internally by Snapchat (based on chat messages and snaps sent in a time window) for such recommendations.

The node set V of our proposed candidate graph $G(V, E)$ is the union of the phone number containing contacts V_p^C (Figure 1a) and the pruned 1-hop friends V_k^F (Figure 1b). We also consider the following variations of graphs that could potentially be used: $V = V_p^C$, $V = V_d^C$, $V = V_p^C \cup V_0^F$, $V = V_d^C \cup V_k^F$ and $V = V_d^C \cup V_0^F$. We define V_d^C as the set of all contacts found in a target user’s contact book that are also on the network; V_0^F as the whole set of 1-hop friends in the network’s friend graph of the previously selected contacts V_p^C or V_d^C , with 0 instead of k in the subscript meaning we don’t prune the selected 1-hop friends based on their number of connections with V_p^C or V_d^C in the network friend graph. We test these candidate graph variations on a dataset of 1.4 million of Snap users using a leave-one-out strategy to simulate new users.

Our utility results are encouraging and can be seen in Table 1. In particular, we find that using the friend candidate graph V_p^C instead of V_d^C slightly improves user engagement performances for a similar structure of the friend candidate graph. Further, including 1-hop

friend nodes (V_0^F or V_k^F) dramatically boosts the number of potential choices for recommended friends and the user engagement, by a factor of 13.25 to 19.5. Although using V_k^F instead of V_0^F gives a much smaller graph, i.e., 40.5% fewer nodes and 48.6% fewer edges, the potential user engagement gain decreases by less than 24.2%, which may be an acceptable cost to pay for raising the privacy bar.

Finally, we A/B tested our system at full scale on the entire Snap network. We found that compared to random candidate selection using $V_p^C \cup V_0^F$, our framework produced a gain of 12.7% new friendships and improved the user engagement up to 6.5%.

5 CONCLUSION

We proposed a framework to recommend friends on social networks while raising the bar on privacy of our users’ friends lists against a brute-force attacker. We found that by utilizing the phone contact book, one can raise the privacy bar and improve the utility of recommendations, as compared with other privacy-preserving alternatives.

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