

Lincheck

Testing concurrency on JVM

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Set-up check

For the workshop you will need:

1. IntelliJ IDEA
2. Cloned [lincheck-workshop](#) project:

```
$ git clone https://git.io/JZp6P
```

Writing concurrent code is pain

Writing concurrent code is pain
... testing it is not much easier!

Lincheck

Lincheck = **L**inearizability **C**hecker (not only linearizability)

<https://github.com/Kotlin/kotlinx-lincheck>

Lincheck

Lincheck = **L**inearizability **C**hecker (not only linearizability)

<https://github.com/Kotlin/kotlinx-lincheck>

1. Generates a random scenario
2. Executes a scenario using either the *stress* or the *model checking strategy*
3. Verifies the results

Counter

```
class Counter {  
    @Volatile var value = 0  
  
    fun addAndGet(delta: Int): Int {  
        value += delta  
        return value  
    }  
}
```

Let's write a test

```
class Counter {  
    @Volatile var value = 0  
  
    fun addAndGet(delta: Int): Int {  
        value += delta  
        return value  
    }  
}
```

```
class CounterTest {  
    private val c = Counter()  
  
    }  
}
```

Initial state

Let's write a test

```
class Counter {  
    @Volatile var value = 0  
  
    fun addAndGet(delta: Int): Int {  
        value += delta  
        return value  
    }  
}
```

```
class CounterTest {  
    private val c = Counter()  
  
    @Operation  
    fun addAndGet(delta: Int) = c.addAndGet(delta)  
}
```

Initial state

Counter operations

Let's write a test

```
class Counter {  
    @Volatile var value = 0  
  
    fun addAndGet(delta: Int): Int {  
        value += delta  
        return value  
    }  
}
```

```
class CounterTest {  
    private val c = Counter()  
  
    @Operation  
    fun addAndGet(delta: Int) = c.addAndGet(delta)  
  
    @Test  
    fun test() = StressOptions()  
        .check(this::class)  
}
```

Initial state

Counter operations

Magic check

Let's write a test

```
class Counter {  
    @Volatile var value = 0  
  
    fun addAndGet(delta: Int): Int {  
        value += delta  
        return value  
    }  
}
```

Run the test!

```
class CounterTest {  
    private val c = Counter()  
  
    @Operation  
    fun addAndGet(delta: Int) = c.addAndGet(delta)  
  
    @Test  
    fun test() = StressOptions()  
        .check(this::class)
```

Initial state

Counter operations

Magic check

Run Counter test

1. Checkout **1.1-counter** branch:

```
$ git checkout 1.1-counter
```

2. CounterTest.kt
3. Run `runStressTest()`, `runModelCheckingTest()`

Failed?

```
class Counter {  
    @Volatile var value = 0  
  
    fun addAndGet(delta: Int): Int {  
        value += delta  
        return value  
    }  
}  
  
class CounterTest {  
    private val c = Counter()  
  
    @Operation  
    fun addAndGet(delta: Int) = c.addAndGet(delta)  
  
    @Test  
    fun test() = StressOptions()  
        .check(this::class)  
}
```

```
java.lang.AssertionError: Invalid  
interleaving found:  
= Invalid execution results: =  
Parallel part:  
| addAndGet(1): 1 | addAndGet(1): 1 |
```

Trace the error

```
class Counter {  
    @Volatile var value = 0  
  
    fun addAndGet(delta: Int): Int {  
        value += delta  
        return value  
    }  
}
```

```
class CounterTest {  
    private val c = Counter()  
  
    @Operation  
    fun addAndGet(delta: Int) = c.addAndGet(delta)  
  
    @Test  
    fun test() = StressOptions()  
        .check(this::class)  
}
```

Stress test

Trace the error

```
class Counter {  
    @Volatile var value = 0  
  
    fun addAndGet(delta: Int): Int {  
        value += delta  
        return value  
    }  
}  
  
class CounterTest {  
    private val c = Counter()  
  
    @Operation  
    fun addAndGet(delta: Int) = c.addAndGet(delta)  
  
    @Test  
    fun test() = ModelCheckingOptions()  
        .check(this::class)  
}
```

Run the test!

Managed strategy

Trace the error

```
org.jetbrains.kotlinx.lincheck.LincheckAssertionError:  
= Invalid execution results =  
Parallel part:  
| addAndGet(1): 1 | addAndGet(2): 2 |  
= The following interleaving leads to the error =  
Parallel part trace:  
| addAndGet(2)  
|   addAndGet(2): 2 at CounterTest.addAndGet(CounterTest.kt:11)  
|     value.READ: 0 at Counter.addAndGet(Counter.kt:7)  
|     switch  
|       addAndGet(1): 1  
|         thread is finished  
|           value.WRITE(2) at Counter.addAndGet(Counter.kt:7)  
|             value.READ: 2 at Counter.addAndGet(Counter.kt:8)  
|             result: 2  
|             thread is finished
```

Trace the error

```
org.jetbrains.kotlinx.lincheck.LincheckAssertionError:
```

```
= Invalid execution results =
```

```
Parallel part:
```

```
| addAndGet(1): 1 | addAndGet(2): 2 |
```

```
= The following interleaving leads to the error =
```

```
Parallel part trace:
```

```
addAndGet(1): 1  
thread is finished
```

```
addAndGet(2)  
addAndGet(2): 2 at CounterTest.addAndGet(CounterTest.kt:11)  
    value.READ: 0 at Counter.addAndGet(Counter.kt:7)  
        switch  
            value.WRITE(2) at Counter.addAndGet(Counter.kt:7)  
                value.READ: 2 at Counter.addAndGet(Counter.kt:8)  
                    result: 2  
                    thread is finished
```

Trace the error

```
org.jetbrains.kotlinx.lincheck.LincheckAssertionError:
```

```
= Invalid execution results =
```

```
Parallel part:
```

```
| addAndGet(1): 1 | addAndGet(2): 2 |
```

```
= The following interleaving leads to the error =
```

```
Parallel part trace:
```

```
addAndGet(1): 1  
thread is finished
```

```
addAndGet(2)  
addAndGet(2): 2 at CounterTest.addAndGet(CounterTest.kt:11)
```

```
value.READ: 0 at Counter.addAndGet(Counter.kt:7)
```

```
switch
```

```
value.WRITE(2) at Counter.addAndGet(Counter.kt:7)
```

```
value.READ: 2 at Counter.addAndGet(Counter.kt:8)
```

```
result: 2
```

```
thread is finished
```



Trace the error

```
org.jetbrains.kotlinx.lincheck.LincheckAssertionError:
```

```
= Invalid execution results =
```

```
Parallel part:
```

```
| addAndGet(1): 1 | addAndGet(2): 2 |
```

```
= The following interleaving leads to the error =
```

```
Parallel part trace:
```

```
addAndGet(1): 1  
thread is finished
```

```
addAndGet(2)  
addAndGet(2): 2 at CounterTest.addAndGet(CounterTest.kt:11)  
value.READ: 0 at Counter.addAndGet(Counter.kt:7)  
switch  
value.WRITE(2) at Counter.addAndGet(Counter.kt:7)  
value.READ: 2 at Counter.addAndGet(Counter.kt:8)  
result: 2  
thread is finished
```

Trace the error

```
org.jetbrains.kotlinx.lincheck.LincheckAssertionError:
```

```
= Invalid execution results =
```

```
Parallel part:
```

```
| addAndGet(1): 1 | addAndGet(2): 2 |
```

```
= The following interleaving leads to the error =
```

```
Parallel part trace:
```

| | |
|---|--|
| <pre>addAndGet(1): 1 thread is finished</pre> | <pre>addAndGet(2) addAndGet(2): 2 at CounterTest.addAndGet(CounterTest.kt:11) value.READ: 0 at Counter.addAndGet(Counter.kt:7) switch value.WRITE(2) at Counter.addAndGet(Counter.kt:7) value.READ: 2 at Counter.addAndGet(Counter.kt:8) result: 2 thread is finished</pre> |
|---|--|



Trace the error

```
org.jetbrains.kotlinx.lincheck.LincheckAssertionError:
```

```
= Invalid execution results =
```

```
Parallel part:
```

```
| addAndGet(1): 1 | addAndGet(2): 2 |
```

```
= The following interleaving leads to the error =
```

```
Parallel part trace:
```

```
addAndGet(1): 1
thread is finished
```



```
addAndGet(2)
addAndGet(2): 2 at CounterTest.addAndGet(CounterTest.kt:11)
value.READ: 0 at Counter.addAndGet(Counter.kt:7)
switch
value.WRITE(2) at Counter.addAndGet(Counter.kt:7)
value.READ: 2 at Counter.addAndGet(Counter.kt:8)
result: 2
thread is finished
```

Trace the error

```
org.jetbrains.kotlinx.lincheck.LincheckAssertionError:  
= Invalid execution results =  
Parallel part:  
| addAndGet(1): 1 | addAndGet(2): 2 |  
= The following interleaving leads to the error =  
Parallel part trace:  
| addAndGet(2)  
|   addAndGet(2): 2 at CounterTest.addAndGet(CounterTest.kt:11)  
|     value.READ: 0 at Counter.addAndGet(Counter.kt:7)  
|     switch  
|       value.WRITE(2) at Counter.addAndGet(Counter.kt:7)  
|       value.READ: 2 at Counter.addAndGet(Counter.kt:8)  
|       result: 2  
|       thread is finished  
addAndGet(1): 1  
thread is finished
```

Correctness

Sequential algorithm



Sequential specification on operations

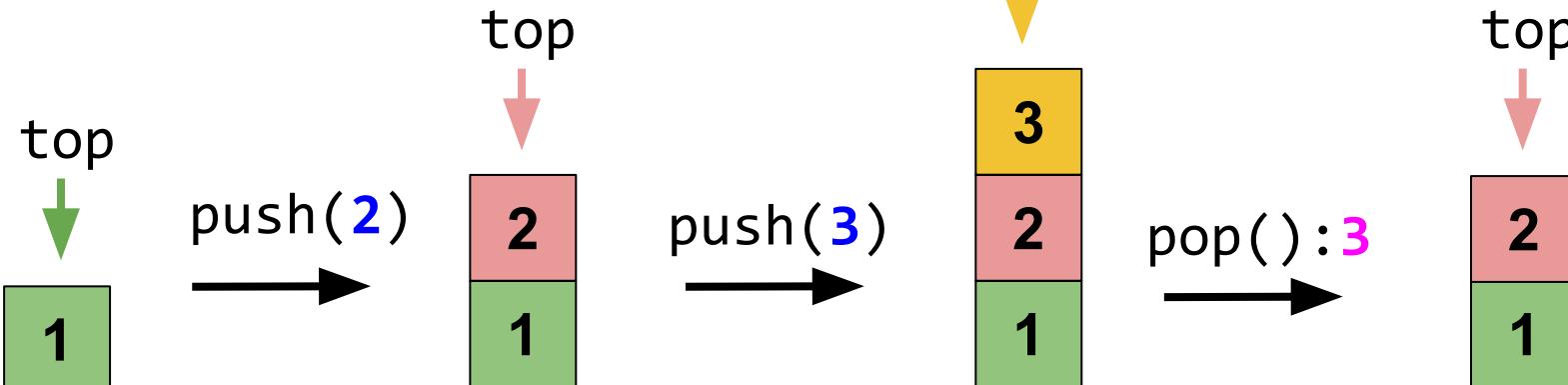
Correctness

Sequential algorithm



Sequential specification on operations

Stack: **LAST IN, FIRST OUT**



Correctness

Sequential algorithm



Sequential specification on operations

Concurrent algorithm



Linearizability (usually)

Correctness

Sequential algorithm



Sequential specification on operations

Concurrent algorithm



Linearizability (usually)

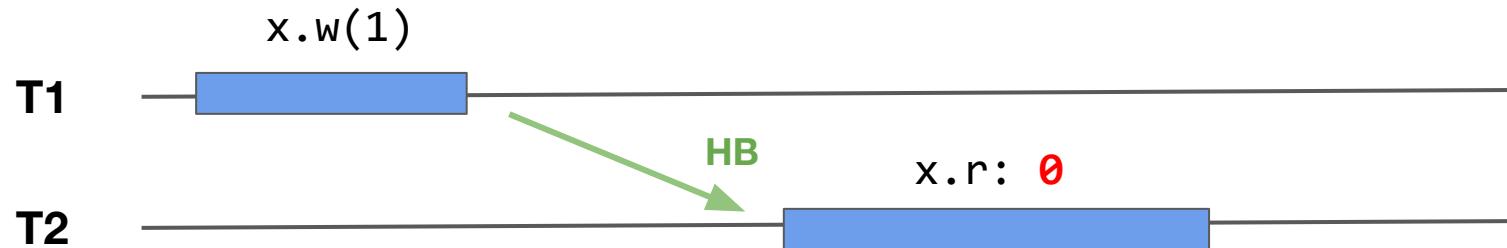
`val c = Counter()`

`c.addAndGet(1): 1`

`c.addAndGet(1): 2`

Execution is **linearizable** iff \exists equivalent sequential execution wrt *happens-before order*

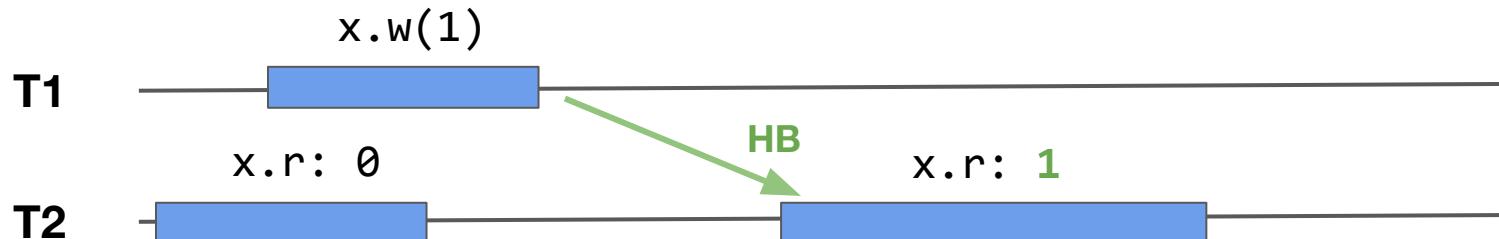
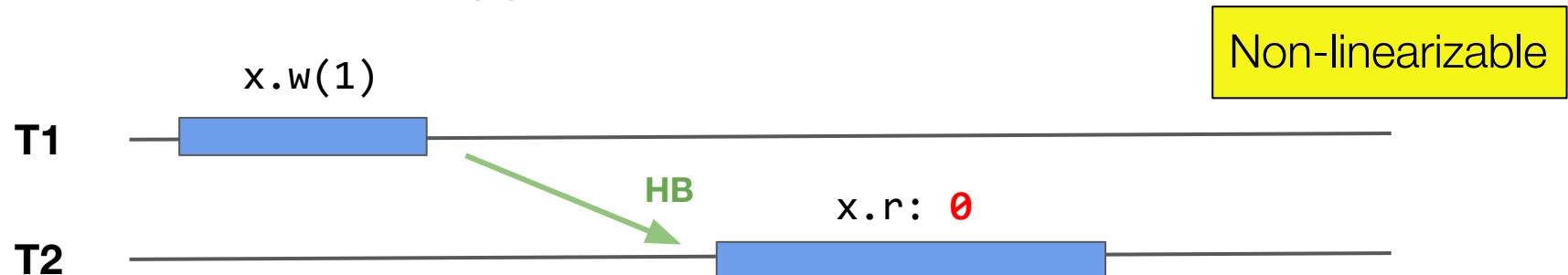
Execution is **linearizable** iff \exists equivalent sequential execution wrt *happens-before* order



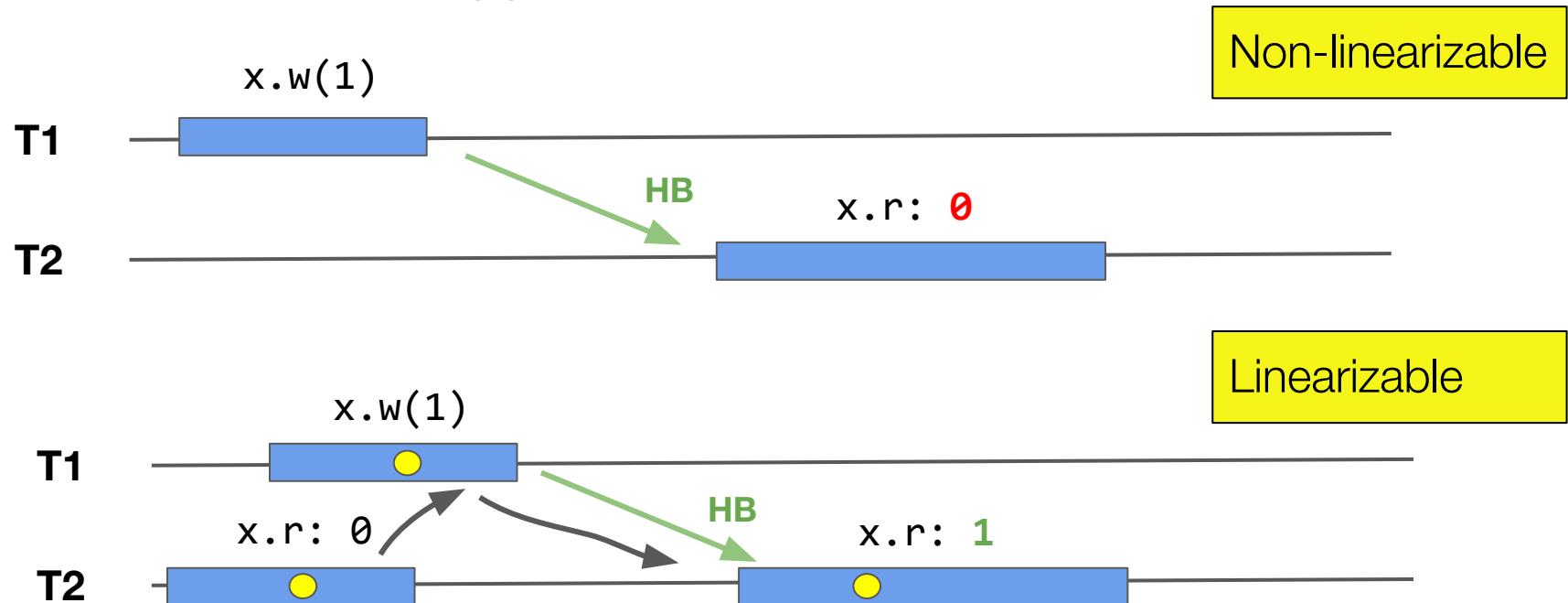
Execution is **linearizable** iff \exists equivalent sequential execution wrt *happens-before* order



Execution is **linearizable** iff \exists equivalent sequential execution wrt *happens-before* order



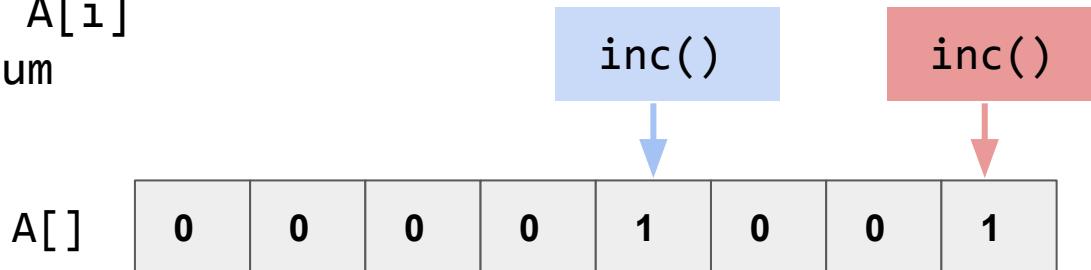
Execution is **linearizable** iff \exists equivalent sequential execution wrt *happens-before* order



Scalable counter

```
fun inc() {  
    i := rand(K)  
    FAA(&A[i], +1)  
}
```

```
fun sum() {  
    sum := 0  
    for (i := 0; i < K; i++)  
        sum += A[i]  
    return sum  
}
```

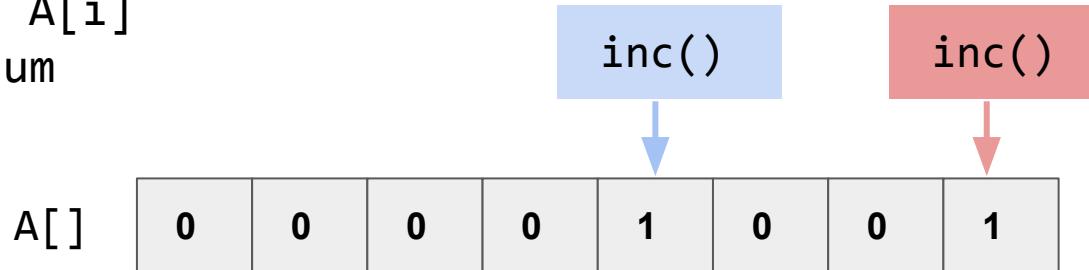


Scalable counter

```
fun inc() {  
    i := rand(K)  
    FAA(&A[i], +1)  
}  
  
fun sum() {  
    sum := 0  
    for (i := 0; i < K; i++)  
        sum += A[i]  
    return sum  
}
```

java.util.concurrent.atomic.LongAdder

- Preferable to AtomicLong when multiple threads update a common sum
- Increases array size under high contention



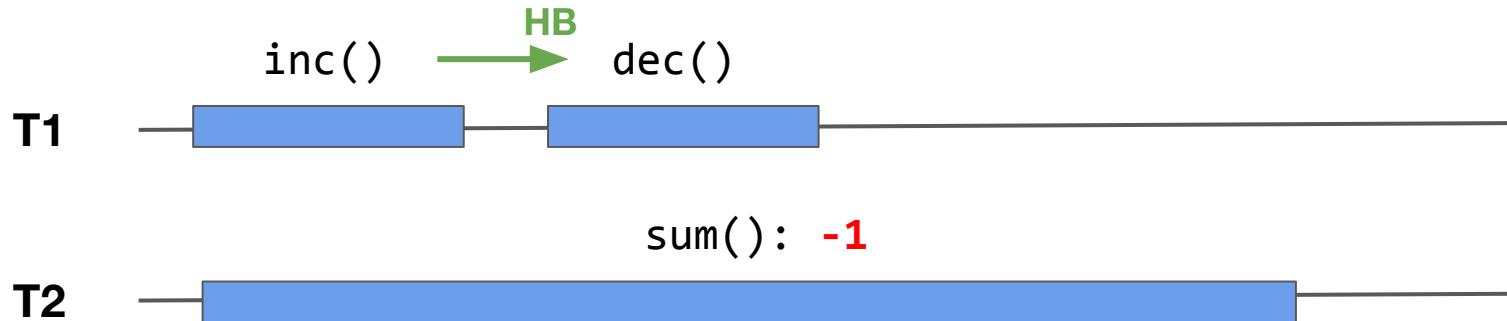
Write a test for LongAdder

1. Checkout **1.2-longAdder** branch:

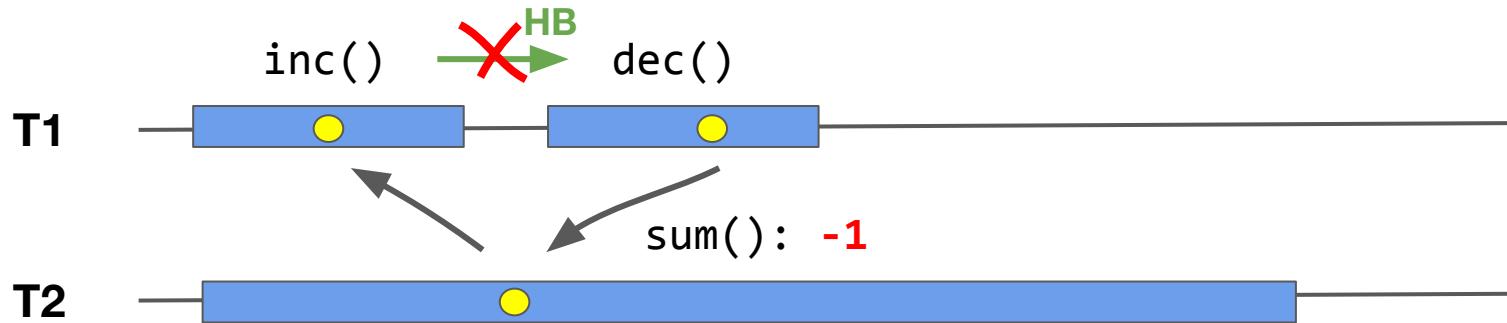
```
$ git checkout 1.2-longAdder
```

2. Complete **TODO()** blocks in `LongAdderTest` class
3. Run `runModelCheckingTest()`

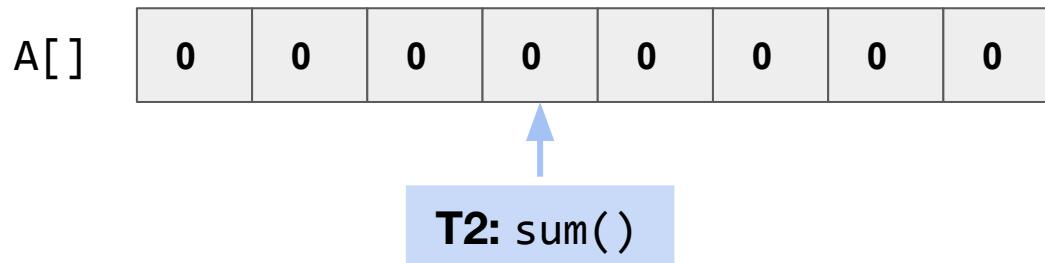
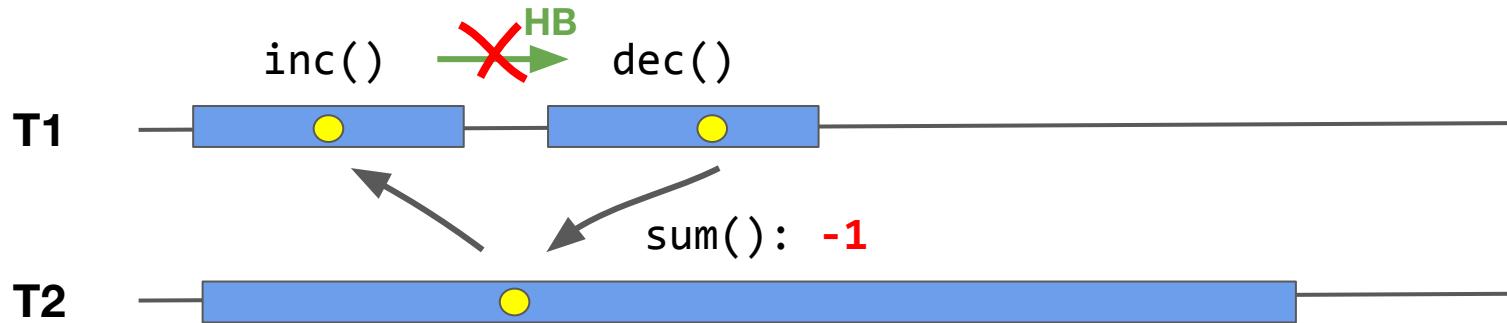
Why LongAdder is not linearizable?



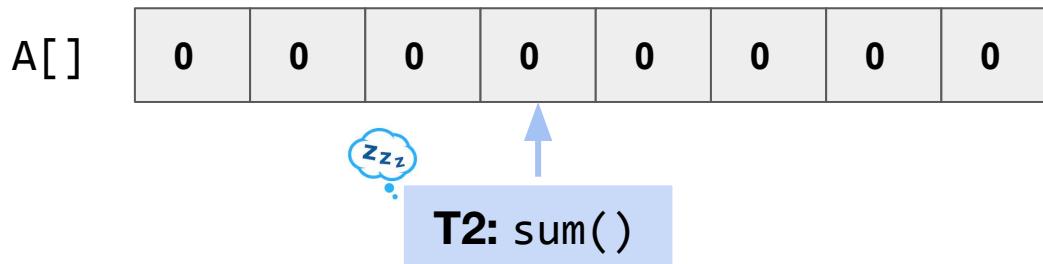
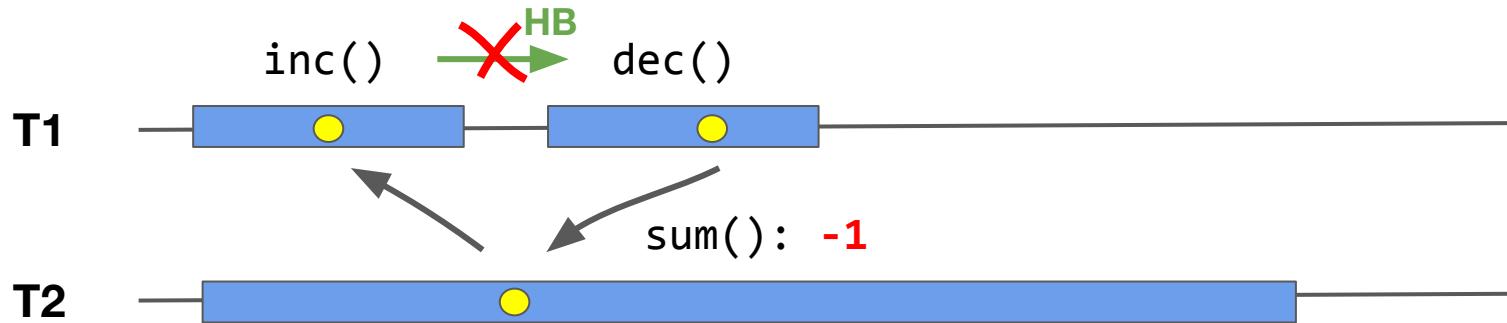
Why LongAdder is not linearizable?



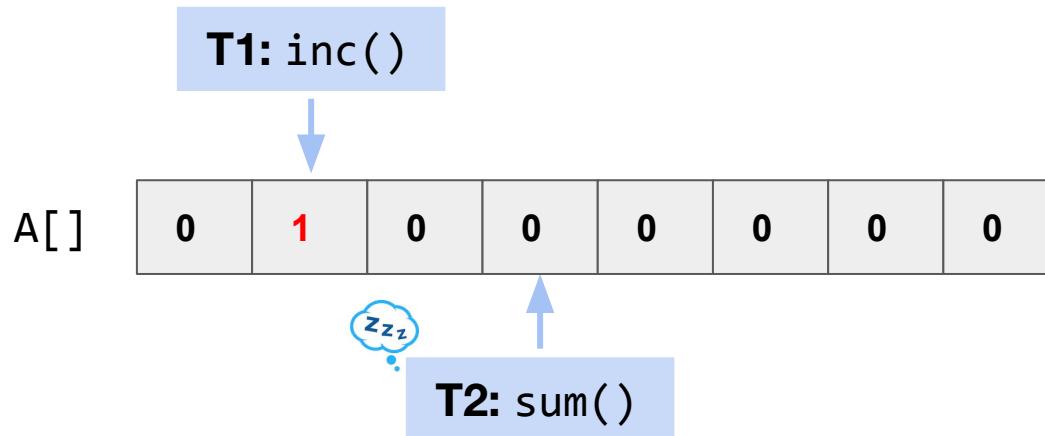
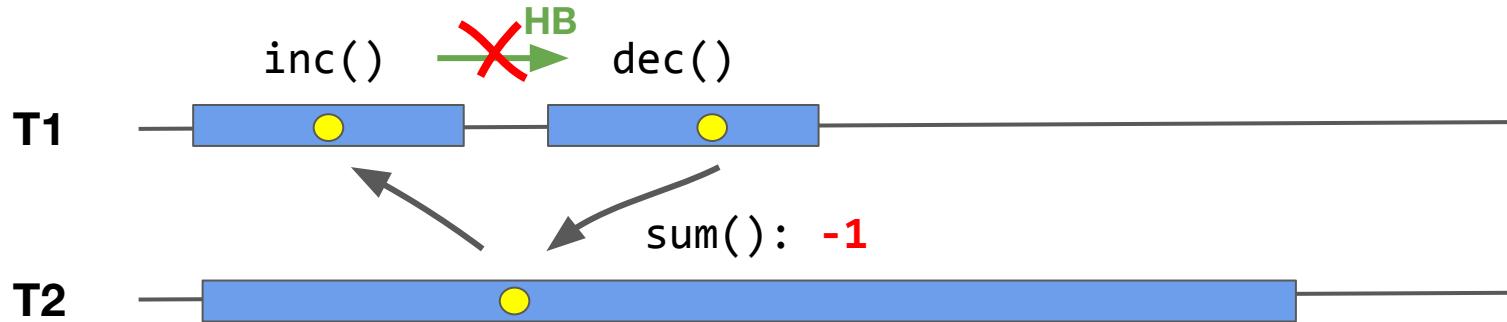
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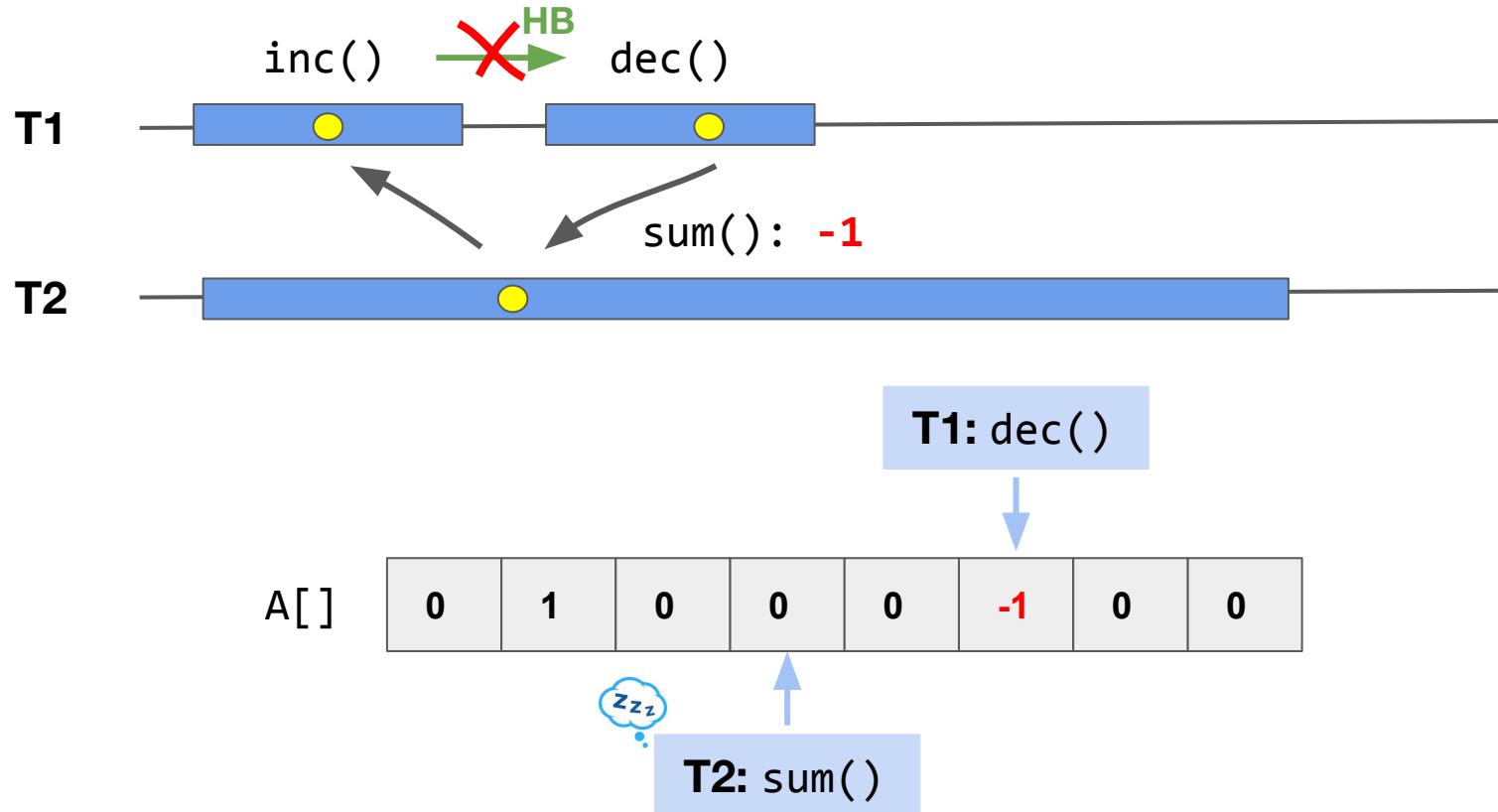
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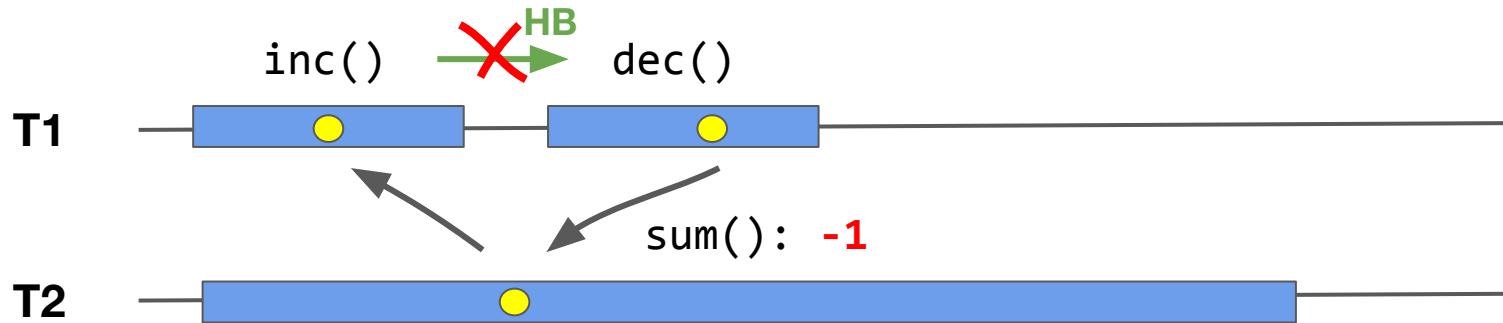
Why LongAdder is not linearizable?



Why LongAdder is not linearizable?



Why LongAdder is not linearizable?



A[]

| | | | | | | | |
|---|---|---|---|---|----|---|---|
| 0 | 1 | 0 | 0 | 0 | -1 | 0 | 0 |
|---|---|---|---|---|----|---|---|

T2: sum(): -1

Testing strategies

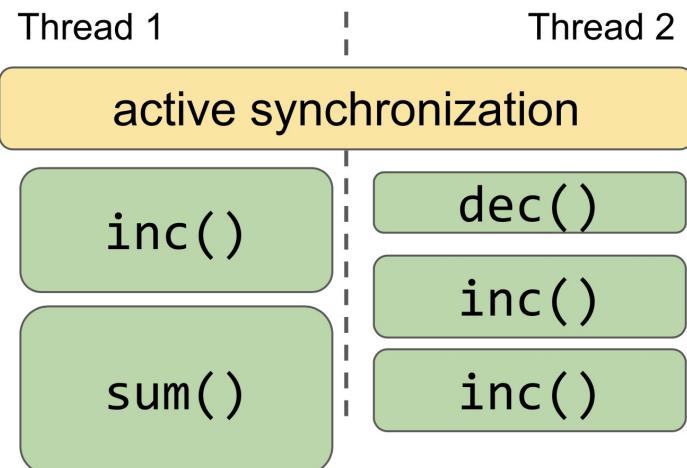
Stress testing

```
class LongAdderTest {  
    ....  
    @Test  
    fun test() = StressOptions()  
        .check(this::class)  
}
```

- Starts real threads
- Actively synchronizes them
- Execute the operation many times

Stress testing

```
class LongAdderTest {  
    ....  
    @Test  
    fun test() = StressOptions()  
        .check(this::class)  
}
```



- Starts real threads
- Actively synchronizes them
- Execute the operation many times

Stress testing

Day 3. June 17



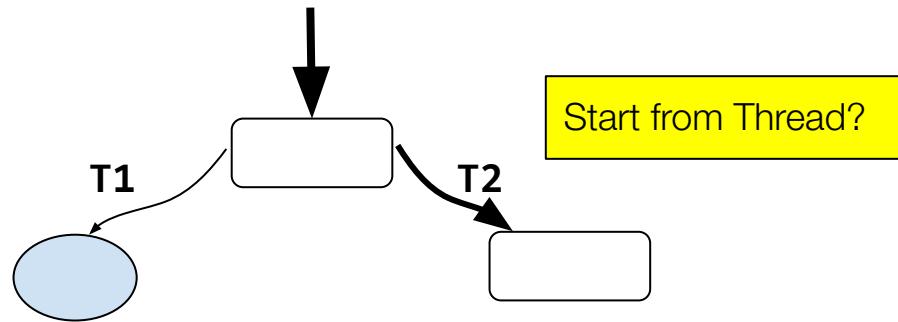
| Time UTC+03:00 | Lecture |
|-------------------|---|
| 18:35 Track 1 | <p>Workshop: Java Concurrency Stress (JCStress) </p> <p>Aleksey Shipilev <i>Red Hat</i></p> <p>#concurrency #java-memory-model #testing #JVM  EN</p> |
| 20:20 Track 1 | <p><u>Workshop: Java Concurrency Stress (JCStress) (part 2)</u> </p> <p>Aleksey Shipilev <i>Red Hat</i></p> <p>#concurrency #java-memory-model #testing #JVM  EN</p> |

Model checking

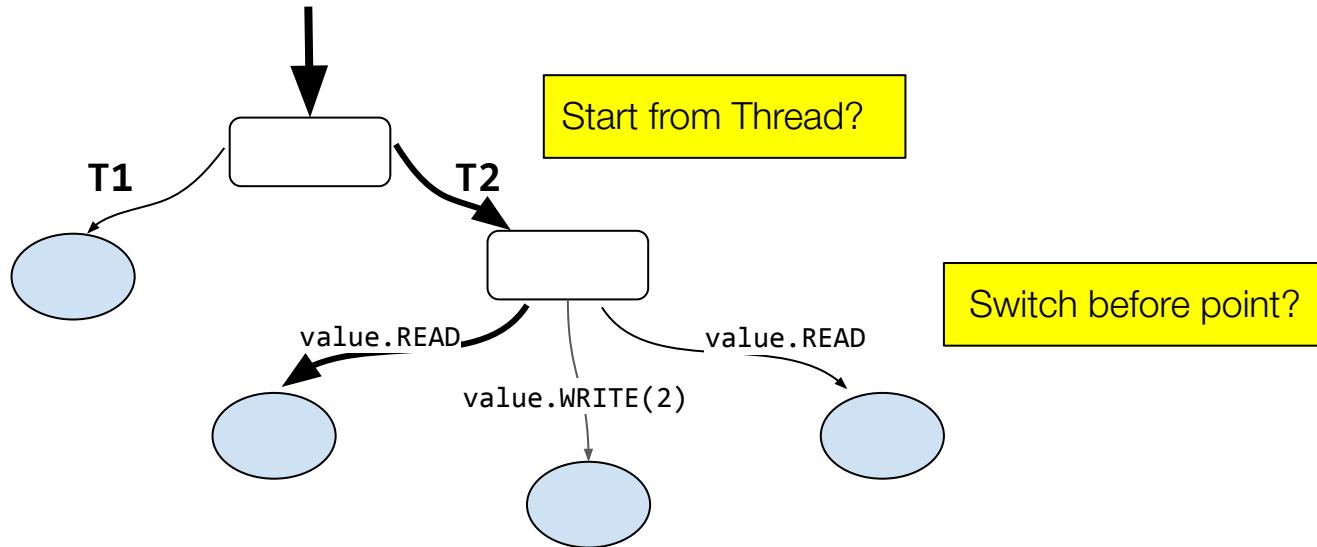
```
class CounterTest {  
    ....  
    @Test  
    fun test() = ModelCheckingOptions()  
        .check(this::class)  
}
```

- Sequential consistency memory model
- Examines many different interleavings within a bounded number of context switches
- Provides a trace

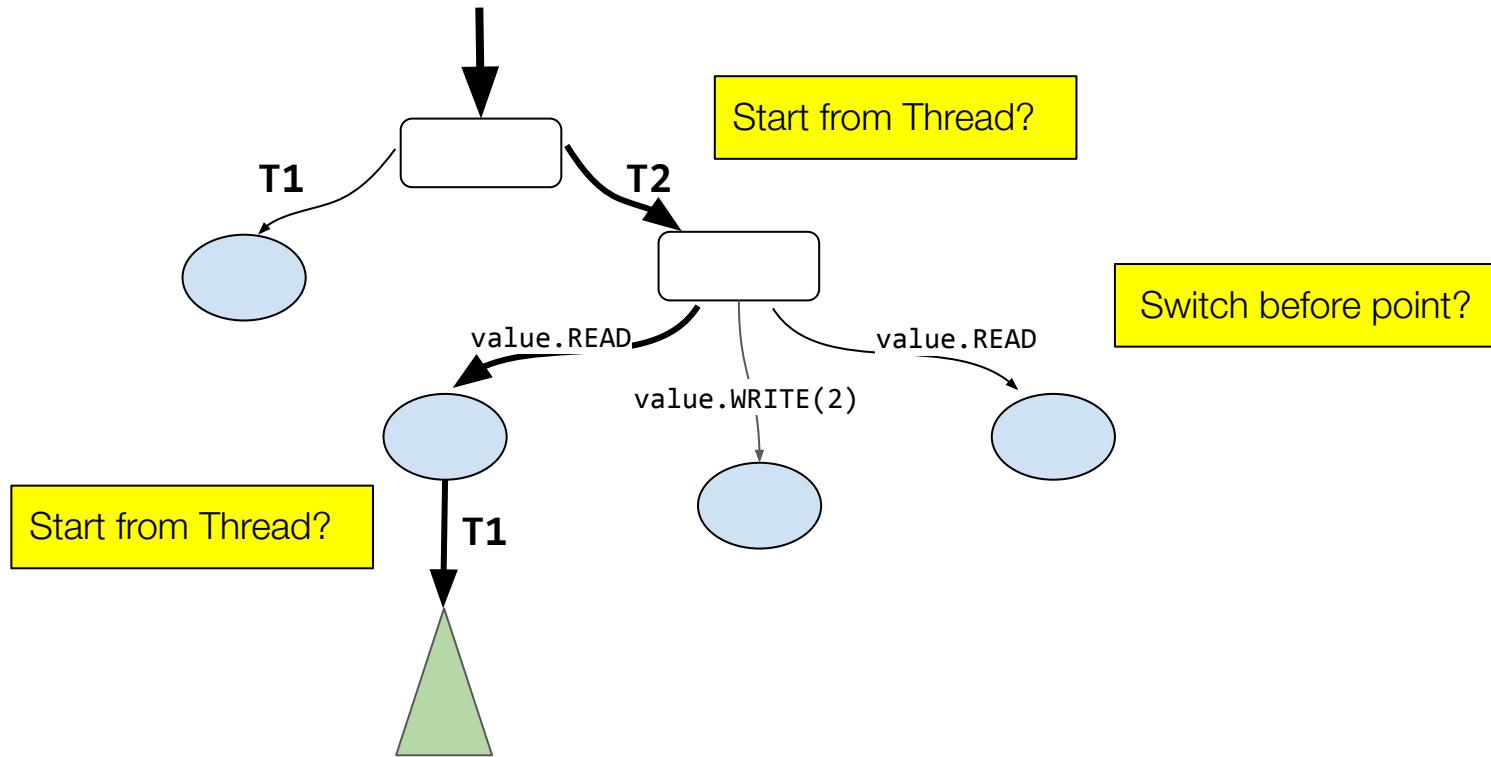
Interleaving tree



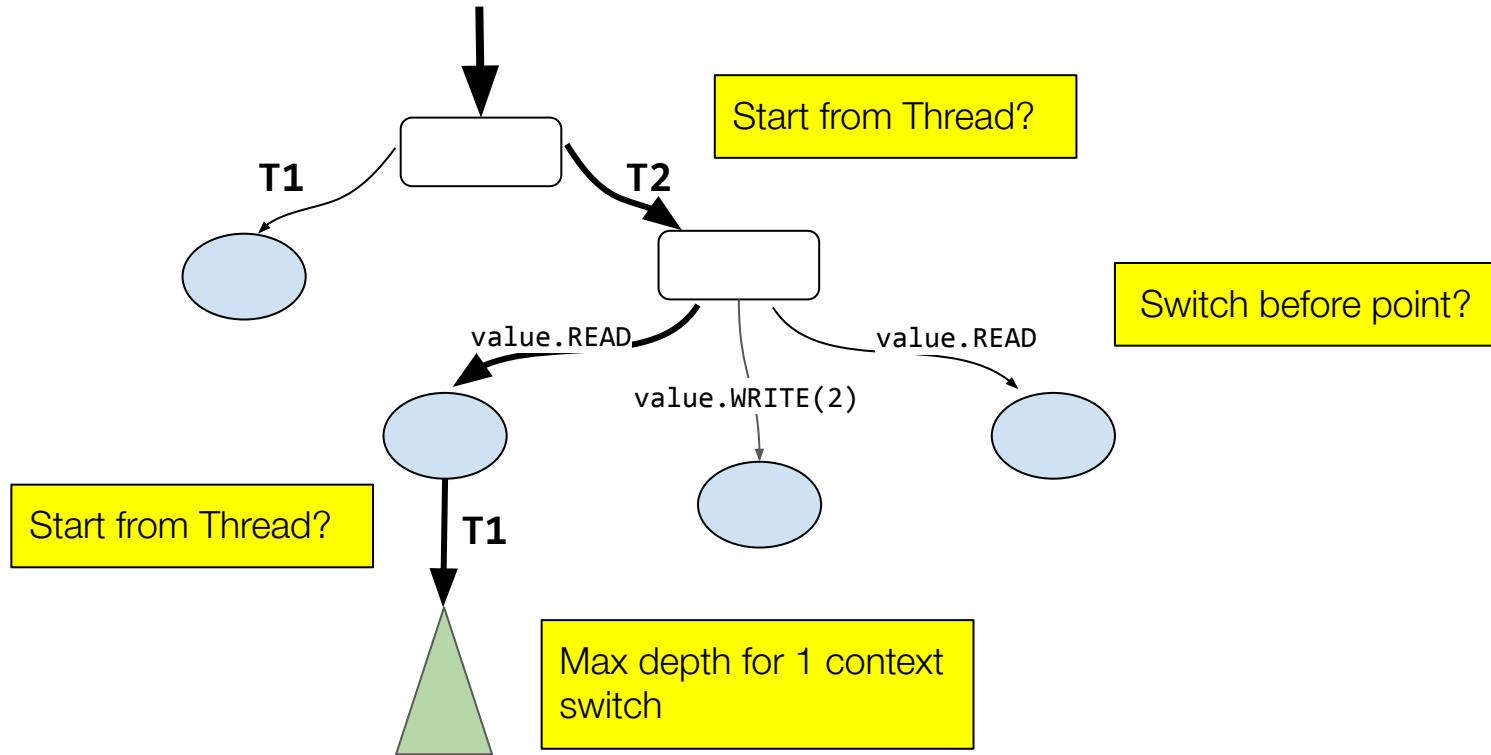
Interleaving tree



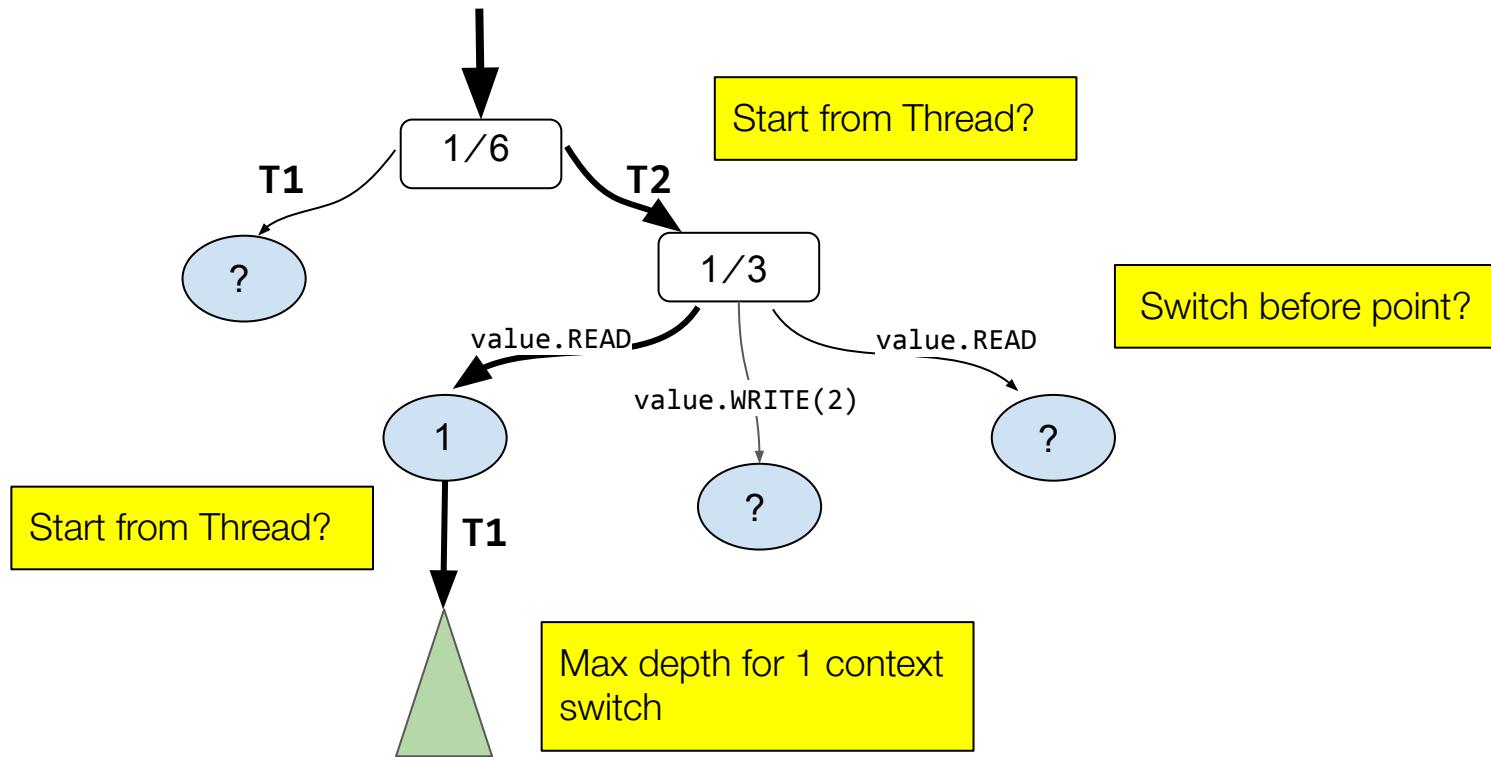
Interleaving tree



Interleaving tree



Interleaving tree



Stress vs Model checking?

Model checking:

- Bugs under sequentially consistent memory model
- Better coverage
- Trace to reproduce the error

Stress testing:

- Low-level effects bugs
- Rare bugs that require many context switches to reproduce (bounds on model checking do not allow that now)

Configurations

Scenario configuration

```
class MyQueueTest {  
    val q = MyQueue<Int>()  
  
    @Operation fun add(x: Int) = q.add(x)  
    @Operation fun poll() = q.poll()  
  
    @Test fun test() = StressOptions()  
        .actorsBefore(2)  
        .threads(2).actorsPerThread(3)  
        .actorsAfter(1)  
        .check(this::class)  
}
```

Scenario configuration

```
class MyQueueTest {  
    val q = MyQueue<Int>()  
  
    @Operation fun add(x: Int) = q.add(x)  
    @Operation fun poll() = q.poll()  
  
    @Test fun test() = StressOptions()  
        .actorsBefore(2)  
        .threads(2).actorsPerThread(3)  
        .actorsAfter(1)  
        .check(this::class)  
}
```

Init part:

[poll(), add(9)]

Parallel part:

| | | |
|--------|--------|--|
| poll() | add(4) | |
| add(3) | add(6) | |
| poll() | poll() | |

Post part:

[add(1)]

Scenario configuration

```
class MyQueueTest {  
    val q = MyQueue<Int>()  
  
    @Operation fun add(x: Int) = q.add(x)  
    @Operation fun poll() = q.poll()  
  
    @Test fun test() = StressOptions()  
        .actorsBefore(2)  
        .threads(2).actorsPerThread(3)  
        .actorsAfter(1)  
        .invocationsPerIteration(10_000) // Run each scenario 10_000 times  
        .iterations(100) // Try 100 different scenarios  
        .check(this::class)  
}
```

Custom scenarios

```
@Test fun test() = StressOptions()  
    .addCustomScenario {  
        parallel {  
            thread {  
                actor(::add, 5)  
                actor(::add, 6)  
            }  
            thread {  
                actor(::poll)  
            }  
        }  
        .check(this::class)
```

Parallel part:

| | |
|--------|--------|
| add(5) | poll() |
| add(6) | |

Stress test example

1. Checkout **2.1-low-level** branch:

```
$ git checkout 2.1-low-level
```

2. WeakMemoryModel.kt
3. Run `test()`

Parameters Generation

```
class MyHashMap {  
    val map = MyHashMap<Int, Int>()  
  
    @Operation  
    fun put(@Param(gen = IntGen::class,  
                  conf = "0:10") key: Int, v: Int) = map.put(key, v)  
  
    @Test fun test() = ...  
}
```

We use parameter generators!

Parameters Generation

```
class MyHashMap {  
    val map = MyHashMap<Int, Int>()  
  
    @Operation  
    fun put(@Param(gen = IntGen::class,  
                  conf = "0:10") key: Int, v: Int) = map.put(key, v)  
  
    @Operation fun get(@Param(gen = IntGen::class,  
                  conf = "0:10") key: Int) = map.get(key)  
  
    @Test fun test() = ...  
}
```

We can share configurations

Parameters Generation

```
@Param(name = "key", gen = IntGen::class, conf = "0:10")
class MyHashMap {
    val map = MyHashMap<Int, Int>()

    @Operation
    fun put(@Param(name = "key") key: Int, v: Int) = map.put(key, v)

    @Operation fun get(@Param(name = "key") key: Int) = map.get(key)

    @Test fun test() = ...
}
```

We can share configurations

Back to LongAdder

1. Checkout **3.1-longAdder** branch:

```
$ git checkout 3.1-longAdder
```

2. Checkout **3.2-jctools** branch

Custom Parameter Generators

```
class RandomIntParameterGenerator(ignoredConf: String)
    : ParameterGenerator<Int> {
    override fun generate() = Random.nextInt()
}
```

Generated values should be Serializable

Constraints

```
class MpscQueueTest {  
    val q = MpscLinkedAtomicQueue<Int>()  
  
    @Operation fun offer(x: Int) = q.offer(x)  
    @Operation fun poll() = q.poll()  
  
    @Test fun test() = ...  
}
```

Parallel part:

| | | |
|-----------|--------------|--|
| add(2) | add(4) | |
| poll(): 2 | poll(): null | |

Constraints

SC queue is incorrect for 2 consumers
by the contract

```
class MpscQueueTest {  
    val q = MpscLinkedAtomicQueue<Int>()  
  
    @Operation fun offer(x: Int) = q.offer(x)  
    @Operation fun poll() = q.poll()  
  
    @Test fun test() = ...  
}
```

Parallel part:

| | | |
|-----------|--------------|--|
| add(2) | add(4) | |
| poll(): 2 | poll(): null | |

Constraints

```
@OpGroupConfig(name = "consumers", nonParallel = true)
class MpscQueueTest {
    val q = MpscLinkedAtomicQueue<Int>()

    @Operation fun offer(x: Int) = q.offer(x)
    @Operation(group = "consumers") fun poll() = q.poll()

    @Test fun test() = ...
}
```

Single-consumer queue

1. Checkout **4.1-constraints** branch

Progress guarantees

An algorithm without explicit synchronization may still be blocking and checking for liveness bugs is a non-trivial task.

Progress guarantees

An algorithm without explicit synchronization may still be blocking and checking for liveness bugs is a non-trivial task.

Lincheck may test the algorithm for **obstruction-freedom** violation.

```
@Test  
fun test() = ModelCheckingOptions()  
    .checkObstructionFreedom()  
    .check(this::class)
```

Progress guarantees

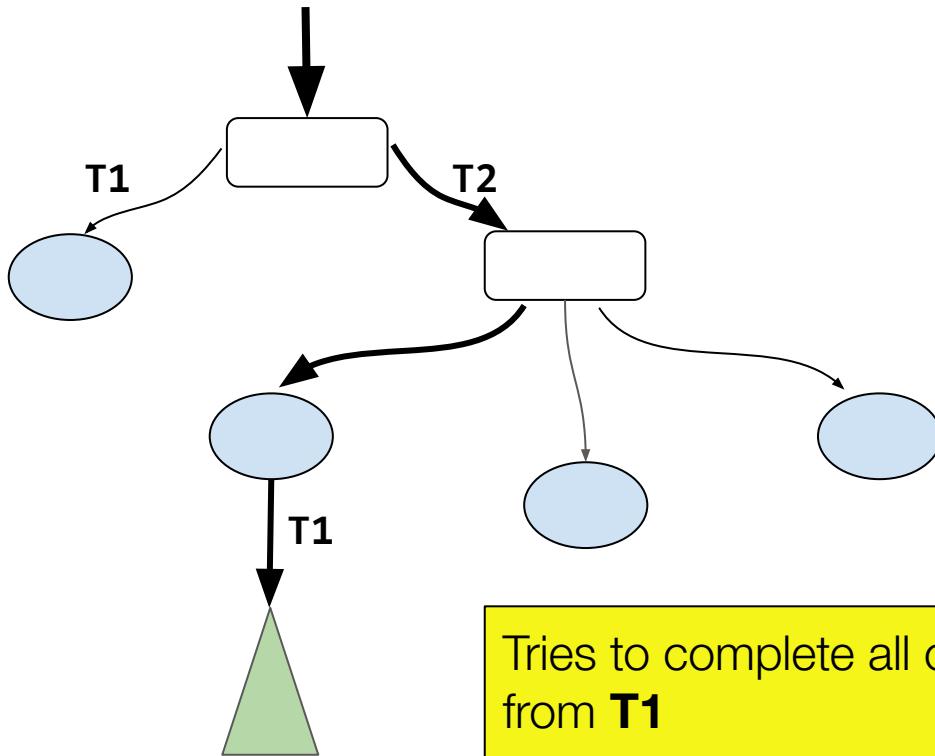
An algorithm without explicit synchronization may still be blocking and checking for liveness bugs is a non-trivial task.

Lincheck may test the algorithm for **obstruction-freedom** violation.

```
@Test  
fun test() = ModelCheckingOptions()  
    .checkObstructionFreedom()  
    .check(this::class)
```

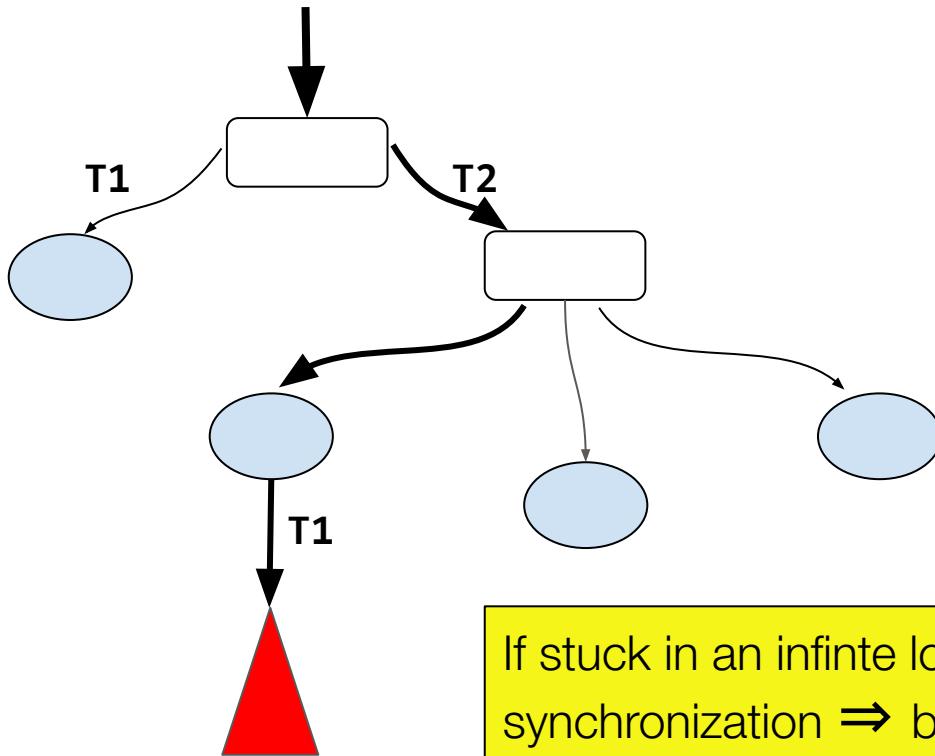
Progress guarantees
are checked in model
checking mode

Progress guarantees



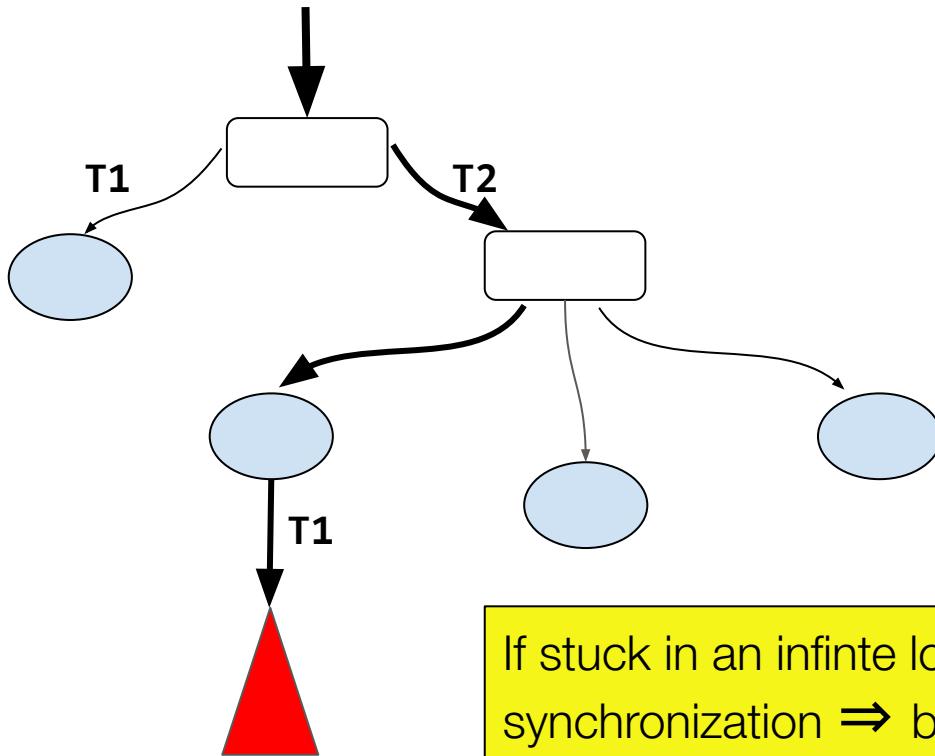
Tries to complete all operations
from **T1**

Progress guarantees



If stuck in an infinite loop or synchronization \Rightarrow blocking

Progress guarantees



Obstruction-freedom: any operation may be completed in a bounded number of steps if all the other processes stop.

If stuck in an infinite loop or synchronization \Rightarrow blocking

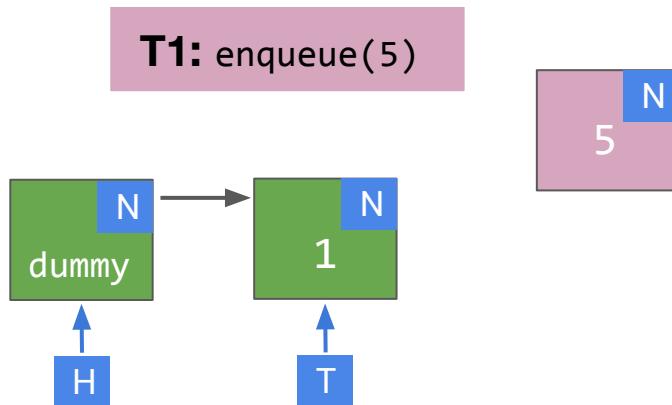
Check whether the algorithm is blocking?

1. Checkout **5.1-dataHolder** branch

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue

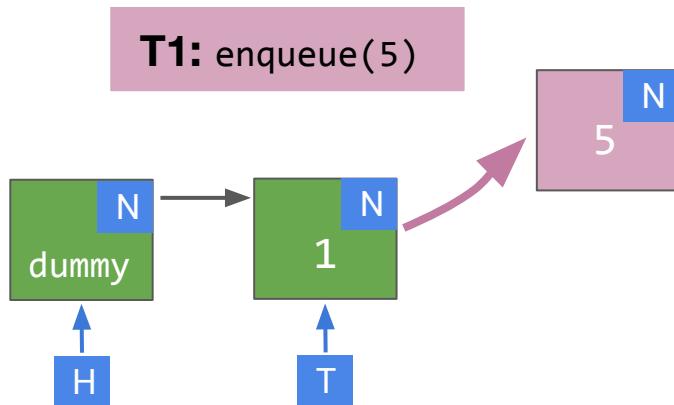


```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        }  
    }  
}
```

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue

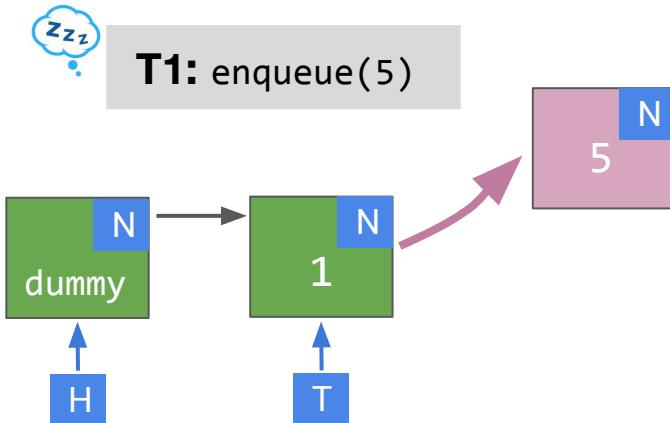


```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        }  
    }  
}
```

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue

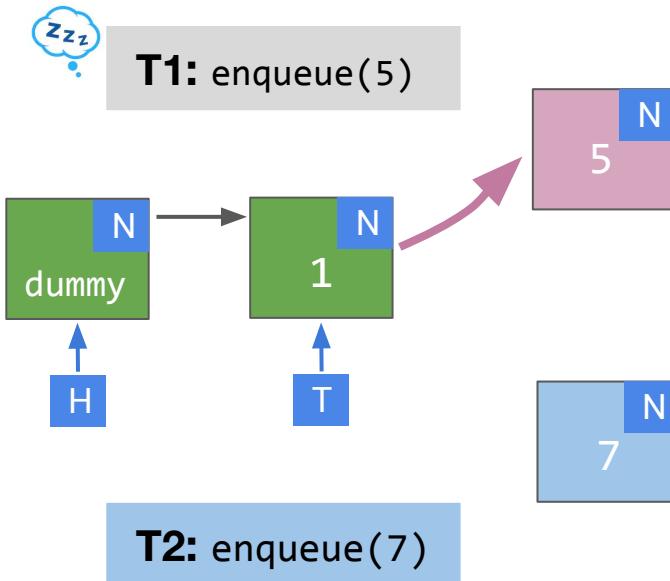


```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        }  
    }  
}
```

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue

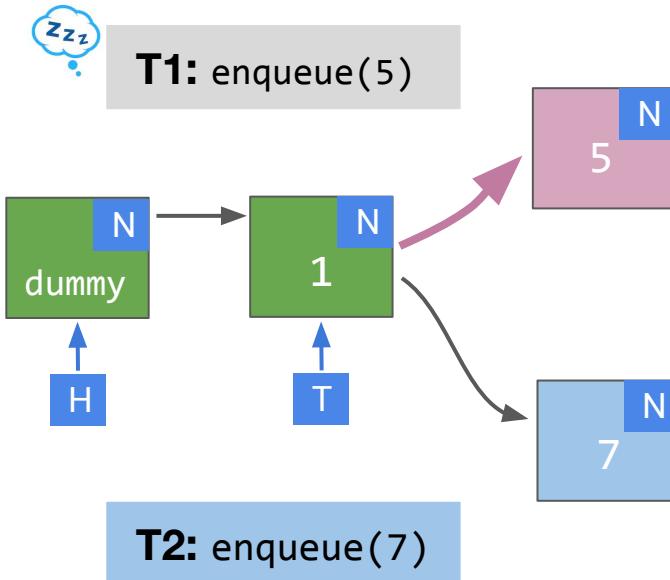


```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        }  
    }  
}
```

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue

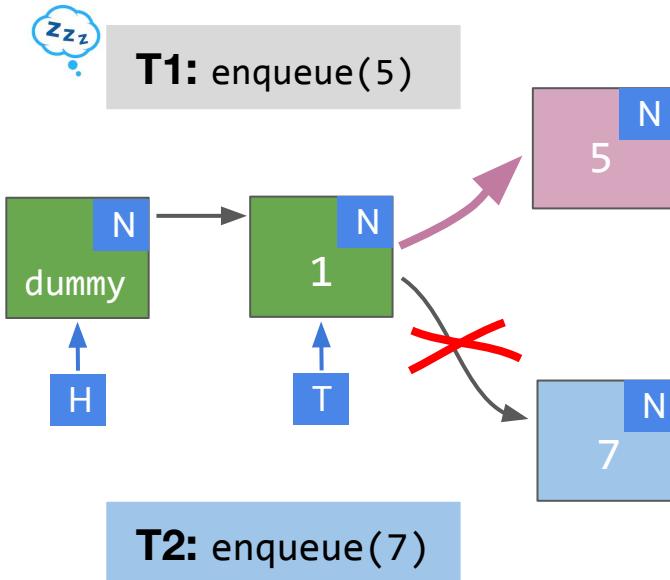


```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        }  
    }  
}
```

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue

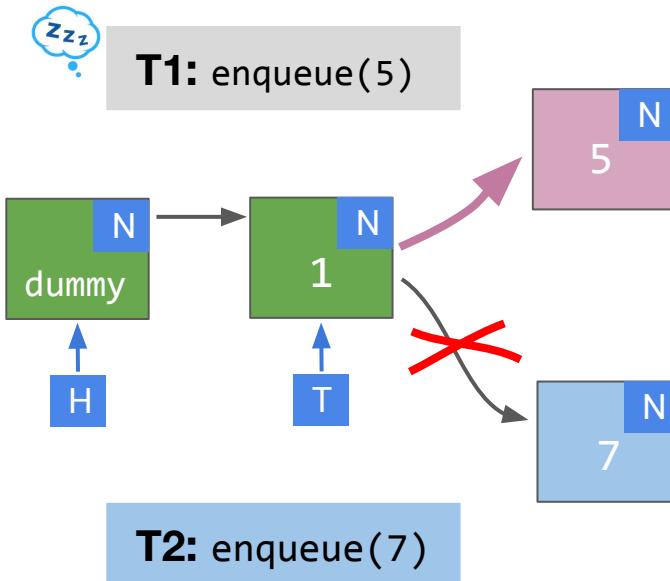


```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        }  
    }  
}
```

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue



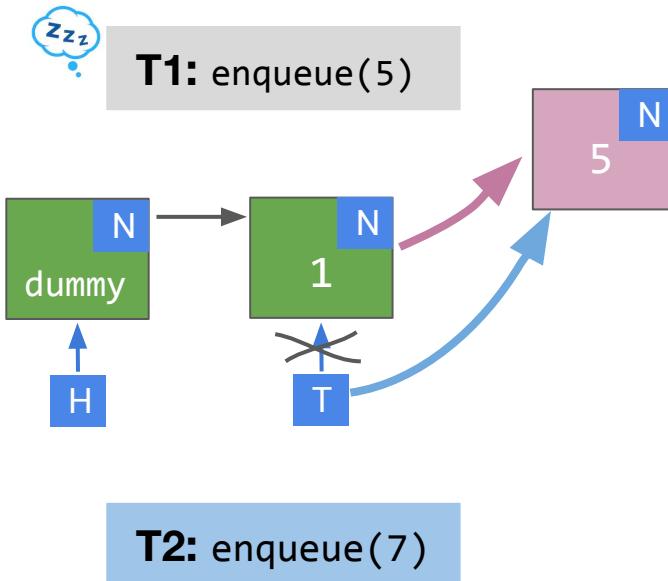
```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        }  
    }  
}
```

T2 can not proceed with
enqueue(7) \Rightarrow blocking
algorithm

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue



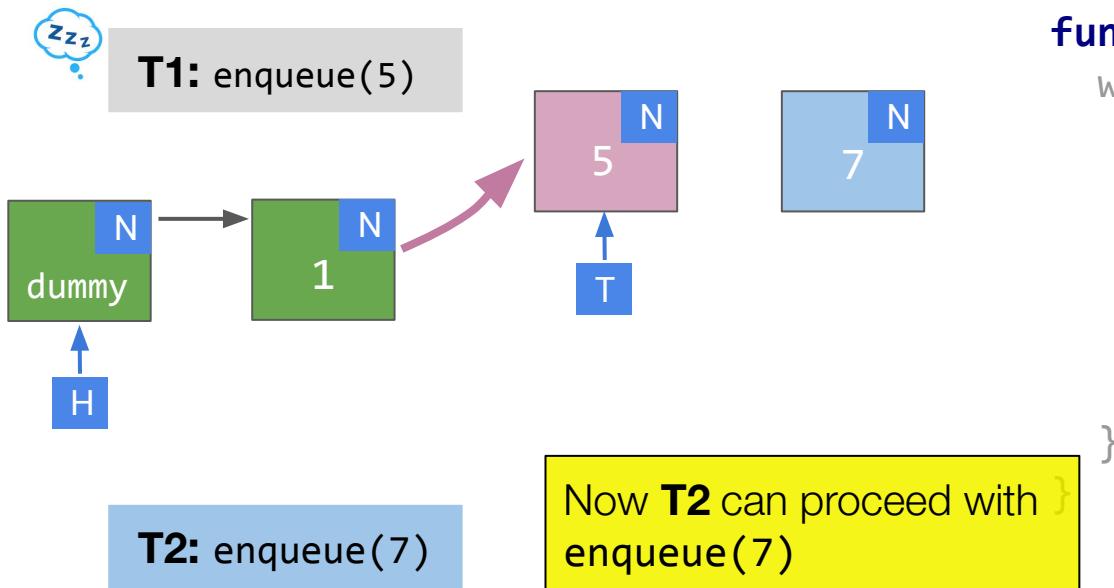
```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        } else {  
            T.CAS(tail, tail.N)  
        }  
    }  
}
```

T2 should help T1 to complete
enqueue(5)

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue

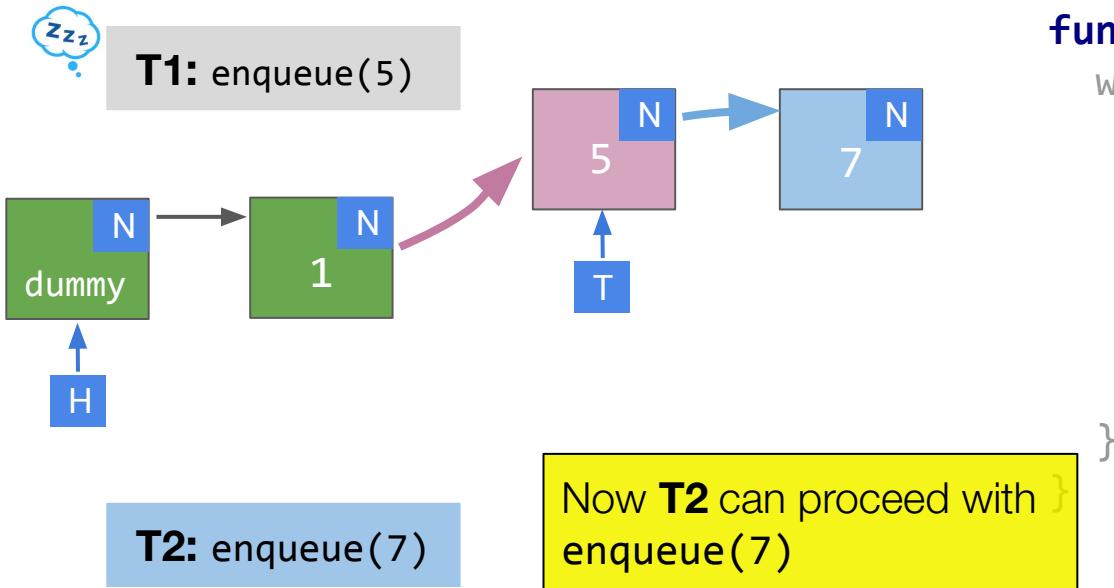


```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        } else {  
            T.CAS(tail, tail.N)  
        }  
    }  
}
```

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue

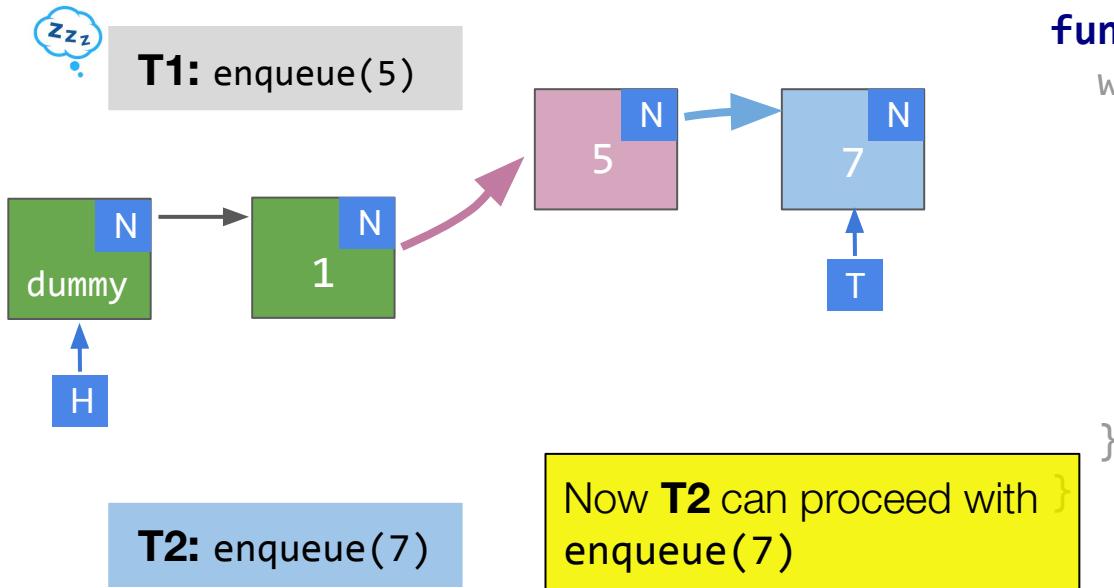


```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        } else {  
            T.CAS(tail, tail.N)  
        }  
    }  
}
```

Helping in lock-free algorithms

Michael-Scott lock-free queue algorithm

the simplest known lock-free queue, j.u.c.ConcurrentLinkedQueue



```
fun enqueue(x) {  
    while (true) {  
        tail := T  
        if (tail.N.CAS(null, x)) {  
            T.CAS(tail, x)  
        } else {  
            T.CAS(tail, tail.N)  
        }  
    }  
}
```

Helping in lock-free algorithms

1. Checkout **5.2-msqueue** branch

Modular testing

If some linearizable data structures are used as building blocks for your algorithm, the number of all possible interleavings may be enormous.

You can treat operations of these data structures atomic \Rightarrow reduce the number of interleavings and speed up testing.

```
@Test
fun test() = ModelCheckingOptions()
    .addGuarantee(
        forClasses(ConcurrentHashMap::class.qualifiedName!!).allMethods().treatAsAtomic())
)
.check(this::class)
```

Modular testing

1. Checkout **6.1-multimap** branch

Verification

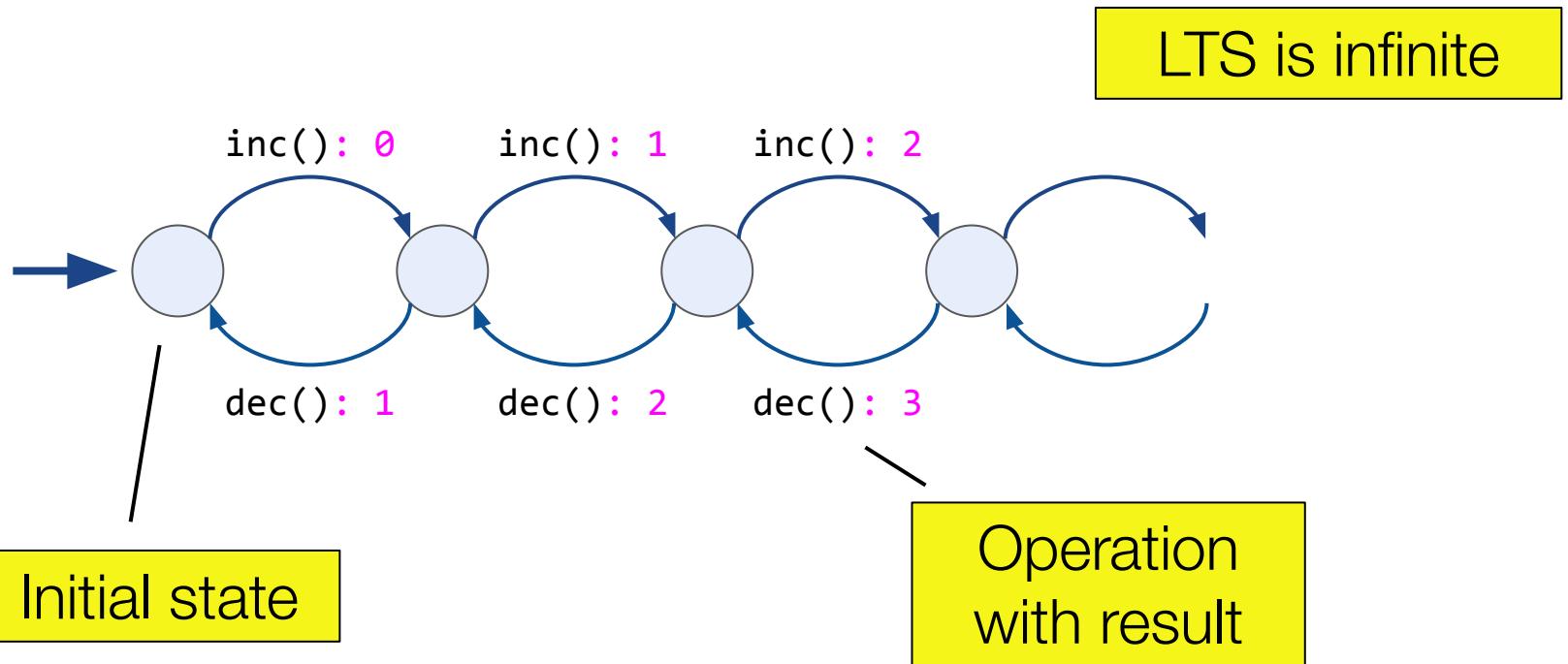
Results Verification

Simplest solution:

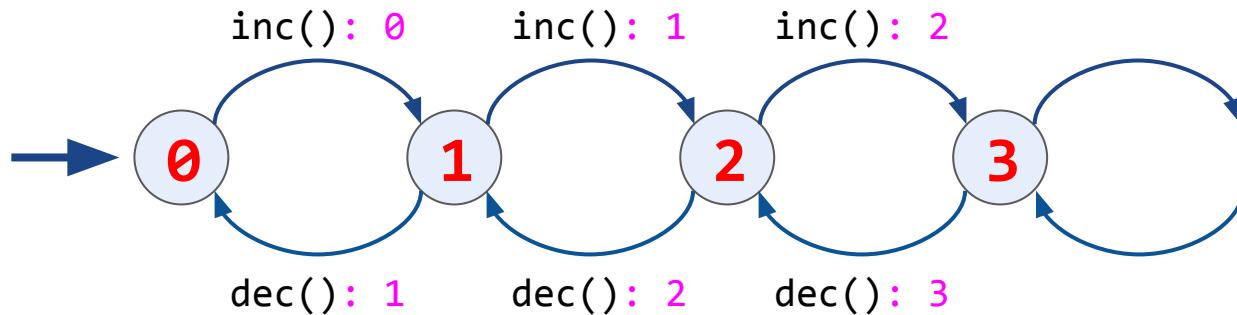
1. Generate all possible sequential histories
2. Check whether one of them produces the same results

Smarter solution: Labeled Transition System (LTS)

LTS (Labeled Transition System)



LTS (Labeled Transition System)



LTS-based verification

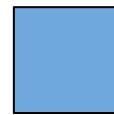
```
val q = MSQueue<Int>()
```

q.add(4)

q.poll(): 9

q.add(9)

q.poll(): 4



LTS-based verification

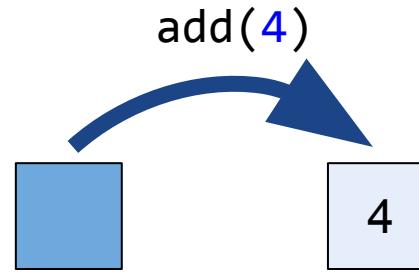
```
val q = MSQueue<Int>()
```

```
q.add(4)
```

```
q.poll(): 9
```

```
q.add(9)
```

```
q.poll(): 4
```

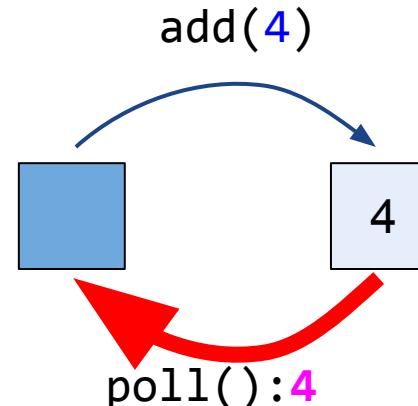


LTS-based verification

```
val q = MSQueue<Int>()
```

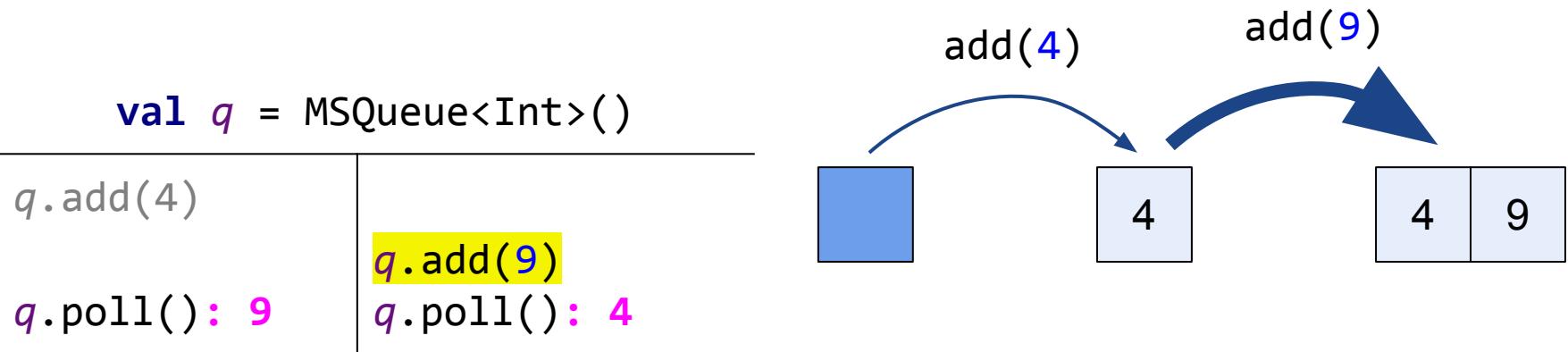
```
q.add(4)  
q.poll(): 9
```

```
q.add(9)  
q.poll(): 4
```



Result is different

LTS-based verification



LTS-based verification

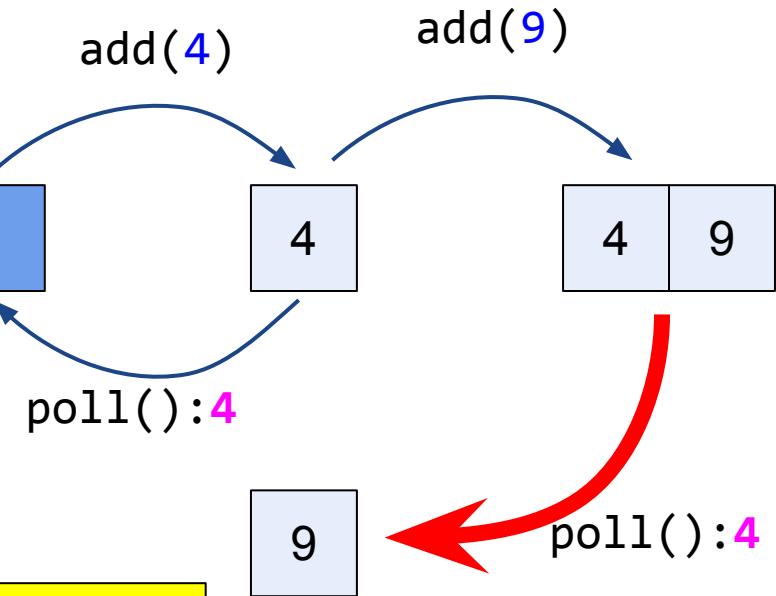
```
val q = MSQueue<Int>()
```

`q.add(4)`

`q.poll(): 9`

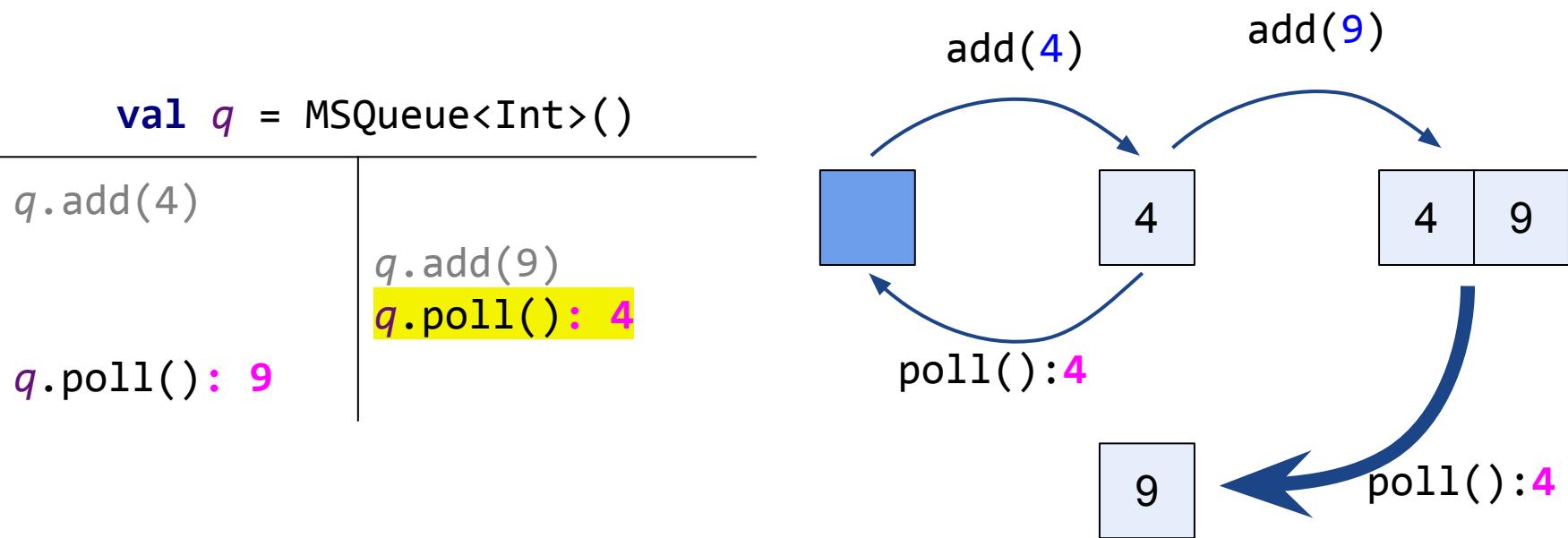
`q.add(9)`

`q.poll(): 4`



Result is different

LTS-based verification



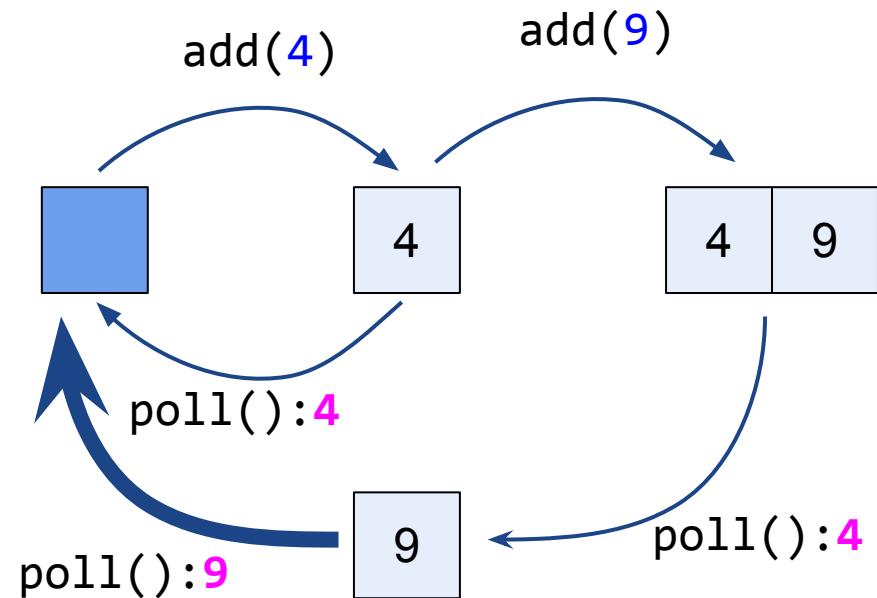
LTS-based verification

`val q = MSQueue<Int>()`

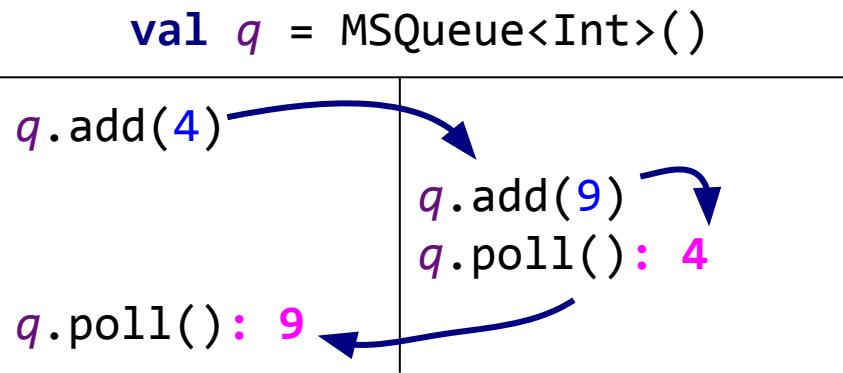
`q.add(4)`

`q.add(9)`
`q.poll(): 4`

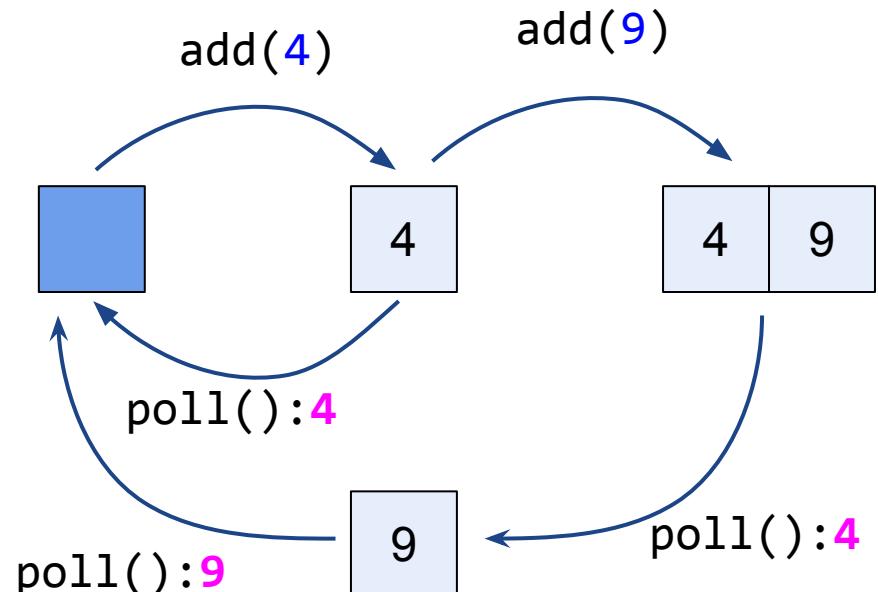
`q.poll(): 9`



LTS-based verification



Path is found \Rightarrow correct



State equivalency

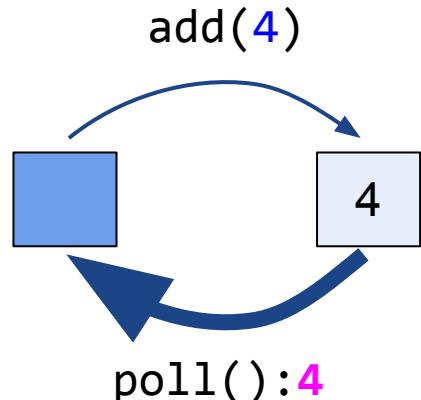
Warning: To make verification faster, you can specify the state equivalence relation on your sequential specification.

Till this moment you got this warning
for all tests

State equivalency

Warning: To make verification faster, you can specify the state equivalence relation on your sequential specification.

We can set state equivalency relation via `equals/hashCode()`:

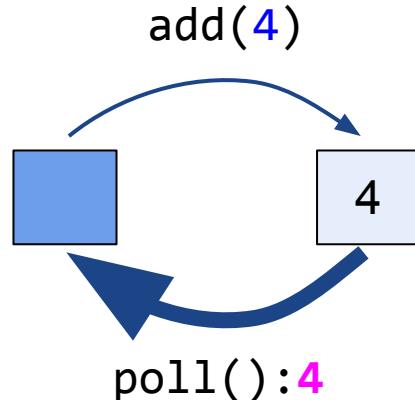


State equivalency

Warning: To make verification faster, you can specify the state equivalence relation on your sequential specification.

We can set state equivalency

relation via `equals/hashCode()`:



```
class MSQueueTest {  
    val q = MSQueue<Int>()  
  
    // Operations here  
  
    override fun equals(other: Any?) = ...  
    override fun hashCode() = ...  
}
```

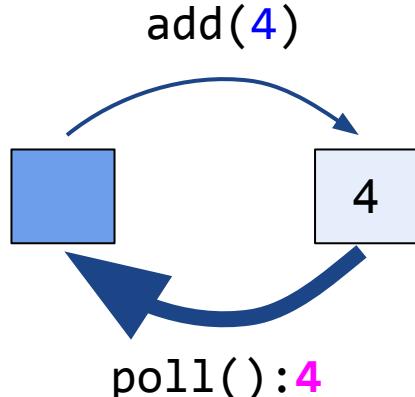
State equivalency

Warning: To make verification faster, you can specify the state equivalence relation on your sequential specification.

We can set state equivalency

relation via `equals/hashCode()`:

```
class MSQueueTest : VerifierState() {  
    val q = MSQueue<Int>()  
  
    // Operations here  
  
    override fun generateState() = q  
}
```



Sequential specification

```
class MySuperFastQueueTest {  
    val q = MySuperFastQueue<Int>()  
  
    @Operation fun offer(x: Int) = q.offer(x)  
    @Operation fun poll() = q.poll()  
  
    @Test fun test() = StressOptions()  
        .sequentialSpecification(SequentialQueue::class.java)  
        .check(this::class)  
}
```

```
class SequentialQueue : VerifierState() {  
    val q = LinkedList<Int>()  
  
    fun offer(x: Int) { q.add(x) }  
    fun poll() = q.remove()  
  
    @Override fun generateState() = q  
}
```

Verification

1. Checkout **7.1-verification** branch

Blocking data structures

Lincheck supports testing blocking data structures implemented via [suspending functions](#) from Kotlin language.

The examples of such data structures from the [Kotlin Coroutines](#) library: [Channel](#), [Mutex](#), [Flow](#)..

See [the coroutines guide](#) for details.

Rendezvous channel

```
class Channel<T> {  
    suspend fun send(x: T)  
    suspend fun receive()  
}
```

```
val c = Channel<Int>()
```

```
c.send(4)
```

```
c.receive() // 4
```

Rendezvous channel

```
class Channel<T> {  
    suspend fun send(x: T)  
    suspend fun receive()  
}
```

```
val c = Channel<Int>()
```

```
c.send(4)
```

```
c.receive() // 4
```

send waits for receive and vice versa

Rendezvous channel

Client 1

```
val task = Task(...)  
tasks.send(task)
```

Client 2

```
val task = Task(...)  
tasks.send(task)
```

Have to wait for send

Worker

```
while(true) {  
    1  val task = tasks.receive()  
    processTask(task)  
}
```

val *tasks* = *Channel*<Task>()

Rendezvous channel

Client 1

```
val task = Task(...)  
tasks.send(task)
```

Client 2

```
val task = Task(...)  
tasks.send(task)
```

Worker



```
while(true) {  
    val task = tasks.receive()  
    processTask(task)  
}
```

val *tasks* = *Channel*<Task>()

Rendezvous channel

Client 1

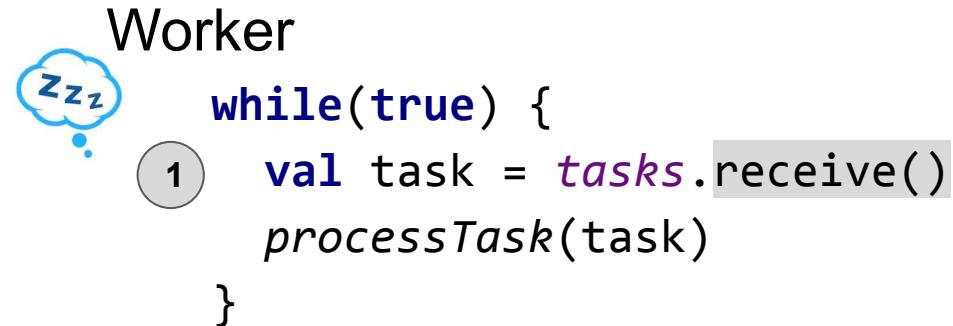
```
val task = Task(...)  
tasks.send(task)
```

Client 2

```
val task = Task(...)  
tasks.send(task)
```

```
val tasks = Channel<Task>()
```

Worker



The diagram shows a worker loop represented by a grey circle containing the number '1'. Above the loop, there is a blue cloud-like shape with three small circles inside, each containing a 'z' character, representing a sleep state. The code within the loop is as follows:

```
while(true) {  
    val task = tasks.receive()  
    processTask(task)  
}
```

Rendezvous channel

Client 1

```
val task = Task(...)
```

2 tasks.send(task)

Rendezvous!

Worker

```
while(true) {
```

1 val task = tasks.receive()
processTask(task)
}

Client 2

```
val task = Task(...)
```

```
tasks.send(task)
```

val tasks = Channel<Task>()

Rendezvous channel

Client 1

```
val task = Task(...)
```

2 tasks.send(task)

Worker

```
while(true) {
```

1

```
    val task = tasks.receive()
```

3

```
    processTask(task)
```

```
}
```

Client 2

```
val task = Task(...)
```

```
tasks.send(task)
```

```
val tasks = Channel<Task>()
```

Rendezvous channel

Client 1

```
val task = Task(...)
```

2 tasks.send(task)

Worker

```
while(true) {
```

1

```
    val task = tasks.receive()
```

3

```
    processTask(task)
```

```
}
```

Client 2

```
val task = Task(...)
```

4 tasks.send(task)

Have to wait for receive

```
val tasks = Channel<Task>()
```

Rendezvous channel

Client 1

```
val task = Task(...)
```

2 tasks.send(task)

Worker

```
while(true) {
```

1

```
    val task = tasks.receive()
```

3

```
    processTask(task)
```

```
}
```

Client 2

```
val task = Task(...)
```

tasks.send(task)

```
val tasks = Channel<Task>()
```



Rendezvous channel

Client 1

```
val task = Task(...)
```

2 tasks.send(task)

Worker

```
while(true) {
```

5

1

```
    val task = tasks.receive()  
    processTask(task)
```

}

Client 2

```
val task = Task(...)
```

4 tasks.send(task)

Rendezvous!

```
val tasks = Channel<Task>()
```

Rendezvous channel

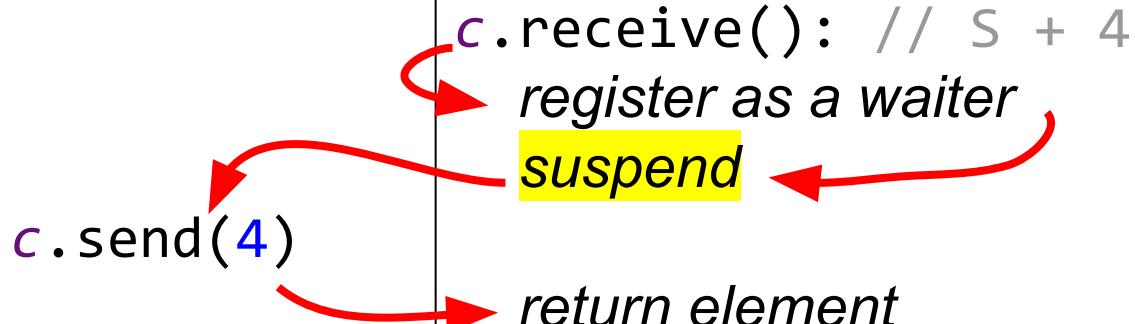
```
val c = Channel<Int>()
```

```
c.send(4)
```

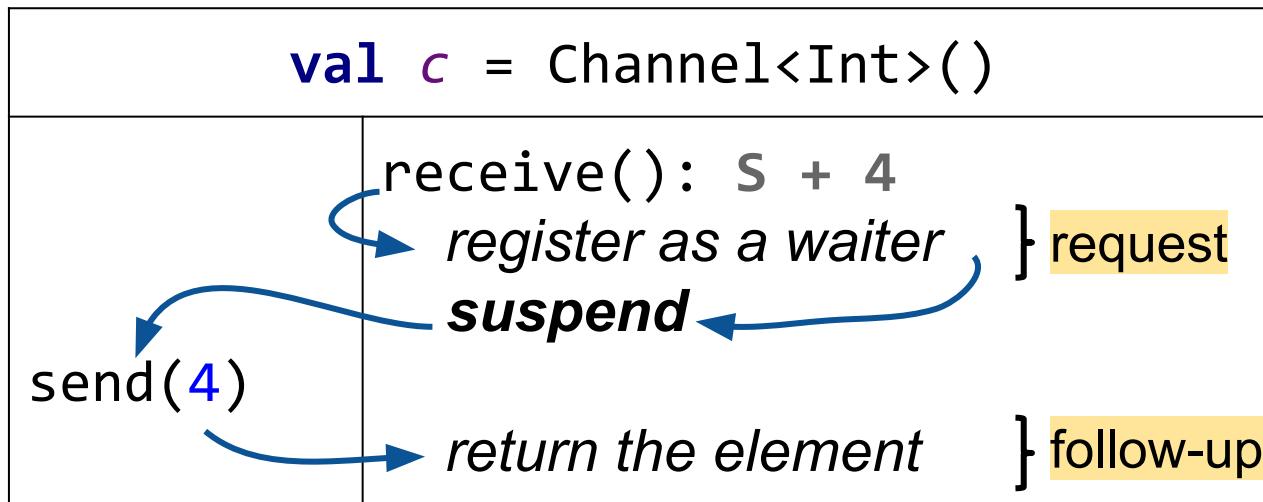
```
c.receive() // S + 4
```

Non-linearizable
because of suspension

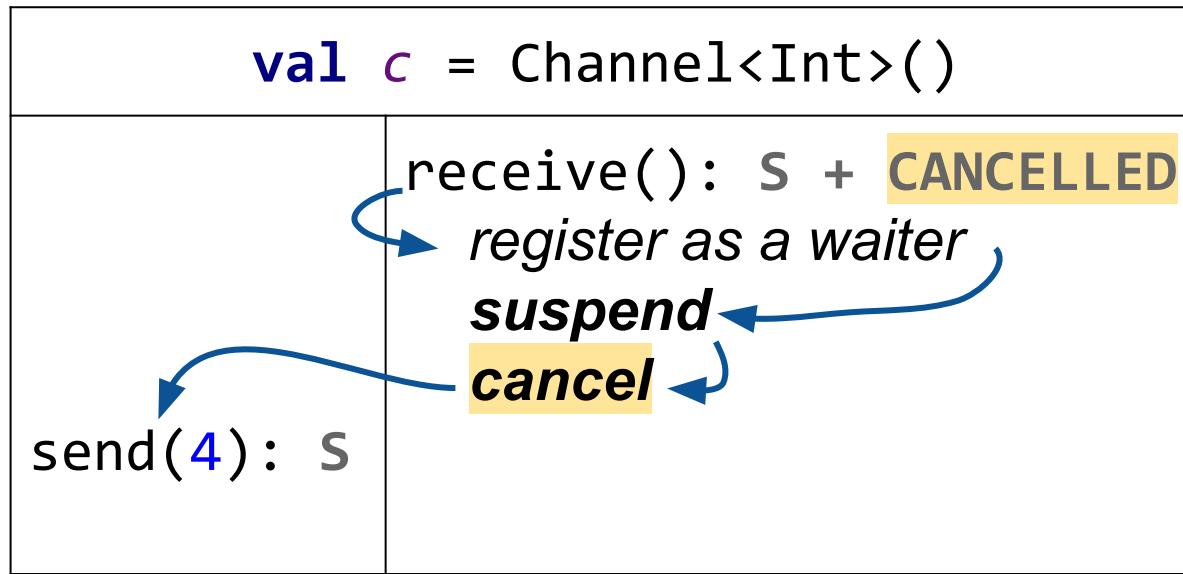
```
val c = Channel<Int>()
```



Dual Data Structures*



Dual Data Structures*



Rendezvous channel test

```
class Channel<T> {  
    suspend fun send(x: T)  
    suspend fun receive()  
}  
  
class ChannelTest {  
    val ch = Channel<Int>()  
  
    @Operation  
    suspend fun send(x: Int) = ch.send(x)  
    @Operation  
    suspend fun receive() = ch.receive()  
  
    @Test fun test() = StressOptions()  
        .check(this::class)  
}
```

Buffered Channels

Client 1

```
val task = Task(...)  
tasks.send(task)
```

Worker

```
while(true) {  
    val task = tasks.receive()  
    processTask(task)  
}
```

Client 2

```
val task = Task(...)  
tasks.send(task)
```

One element can be sent
without suspension

```
val tasks = Channel<Task>(capacity = 1)
```

Buffered Channels

Client 1

```
val task = Task(...)
```

```
1 tasks.send(task)
```

Worker

```
while(true) {  
    val task = tasks.receive()  
    processTask(task)  
}
```

Client 2

```
val task = Task(...)
```

```
tasks.send(task)
```

Does not suspend!

```
val tasks = Channel<Task>(capacity = 1)
```

Buffered Channels

Client 1

```
val task = Task(...)
```

1 tasks.send(task)

Worker

```
while(true) {  
    val task = tasks.receive()  
    processTask(task)}
```

Client 2

```
val task = Task(...)
```

2 tasks.send(task)

The buffer is full, suspends

```
val tasks = Channel<Task>(capacity = 1)
```

Buffered Channels

Client 1

```
val task = Task(...)
```

1 tasks.send(task)

Client 2

```
val task = Task(...)
```

2 tasks.send(task) —————

Worker

```
while(true) {  
    val task = tasks.receive()  
    processTask(task)  
}
```

Receives the buffered element,
resumes the 2nd client,
and moves its task to the buffer

```
val tasks = Channel<Task>(capacity = 1)
```

Buffered Channels

Client 1

```
val task = Task(...)
```

1 tasks.send(task)

Client 2

```
val task = Task(...)
```

2 tasks.send(task)

Worker

```
while(true) {
```

3 val task = tasks.receive()
processTask(task)

4 }

}

Retrieves the 2nd task,
no waiters to resume

```
val tasks = Channel<Task>(capacity = 1)
```

Blocking data structures

1. Checkout **8.1-rendezvous-channel** branch
2. Checkout **8.2-buffeed-channel** branch

For deep dive into verification of blocking data structures: [“LinCheck: Testing concurrent data structures in Java”](#) by Nikita Koval @ Hydra 2019

St. Petersburg
July 11-12 2019

Lazy Dual Data Structures LTS creation

defined via equals/hashcode

States are equal iff $\exists f : \mathbb{N} \rightarrow \mathbb{N}$ that

1. externally observable states
2. $st-s$ wrt rf on tickets (as lists)
3. $rt-s$ wrt rf on tickets (as sets) are equal

maintained by Lin-Check

The diagram illustrates the construction of an LTS for a lazy dual data structure. It shows states represented by boxes and transitions between them. States include a single ticket (e.g., 4), a list of tickets (e.g., [4, 4]), and a set of tickets (e.g., {2 → 1}). Transitions are labeled with actions like `c.send(0, 4)` and `c.receive(0)`, along with their resulting states and ticket relations (rf). A yellow box at the top right defines state equality via `equals`/`hashcode`. A red box contains the three conditions for state equality. A yellow box at the bottom right indicates that the LTS is maintained by Lin-Check.

Questions?