Valhalla is coming

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Java Platform Group Oracle June, 2020



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Who am I?

- Java/JVM Performance Engineer at Oracle, @since 2010
- Java/JVM Performance Engineer, @since 2005
- Java/JVM Engineer, @since 1996

Preface



3 Key Principles of Java Evolution

- Compatibility
- Compatibility
- Compatibility

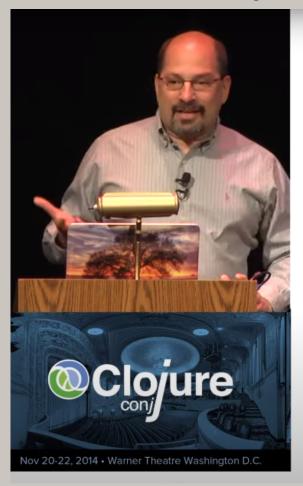


- If we don't know how to do it right, we won't do it (now).
- It's better to leave out something good than to put in something bad. (Because compatibility is forever)
- Less is more The first Java "buzzword" is SIMPLE

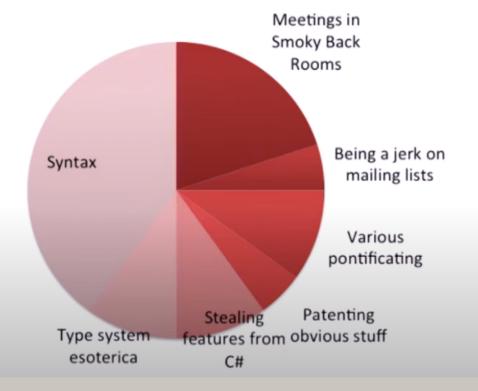
(c) Brian Goetz

Java Evolution

Brian Goetz: "Stewardship: the Sobering Parts"



What people think I do (naïve version)

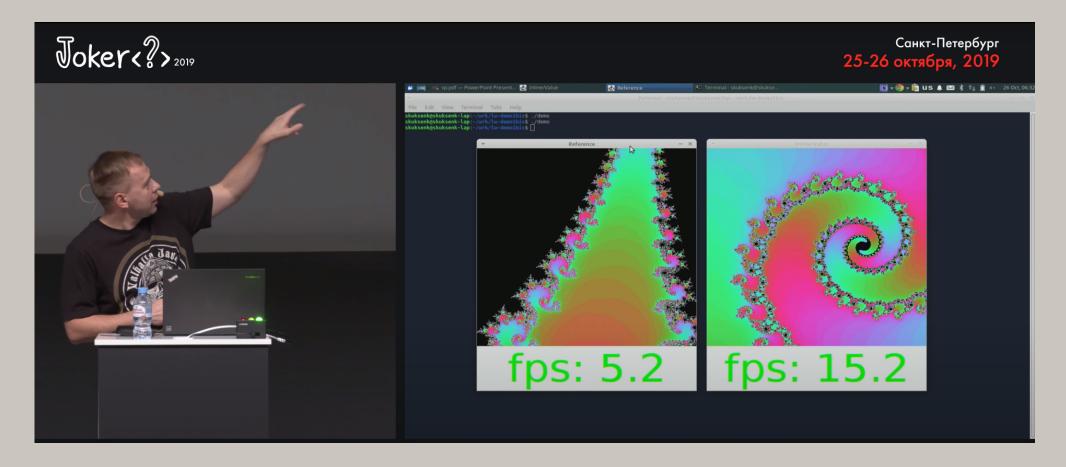


What is Valhalla?



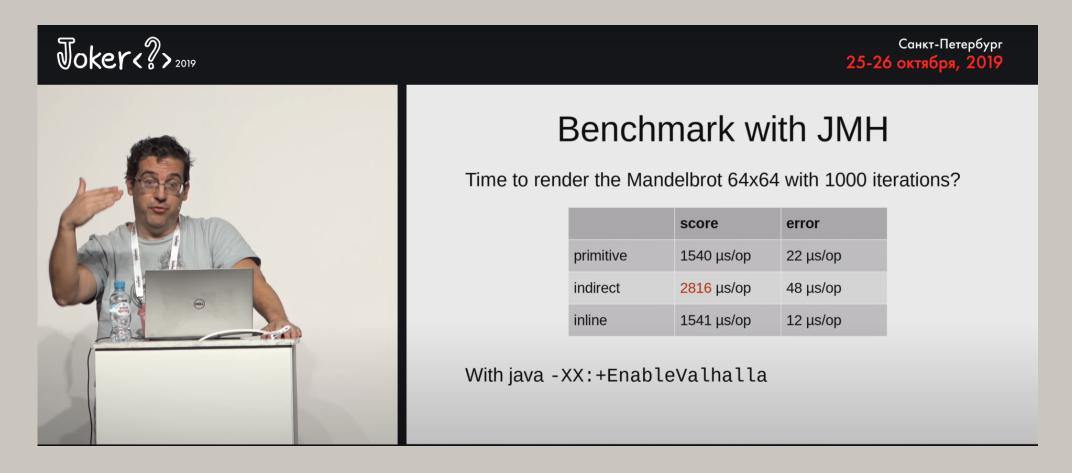
What is Valhalla?

"Does Java need "inline" types?"



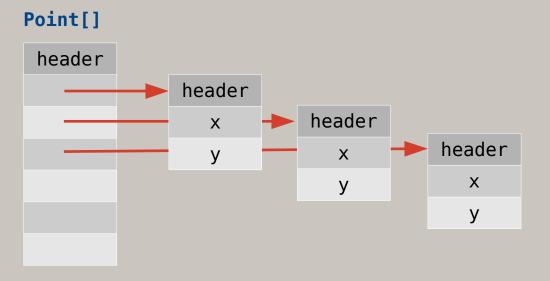
What is Valhalla?

Remi Forax: "The sinuous path toward Valhalla"



Inefficiency: references are everywhere

class Point {...}



Solution: Inline types

```
inline class Point {...}
```

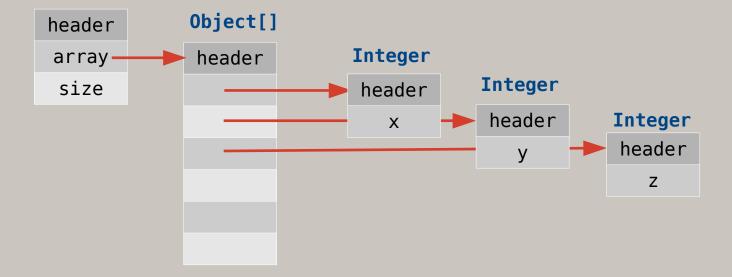
Point[]

header

x
y
x
y
x
y
y

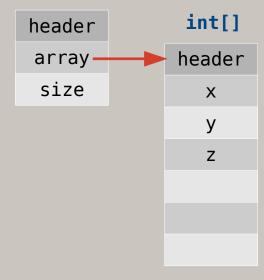
Inefficiency: boxes are everywhere

ArrayList<Integer>



Solution: Specialized generics

ArrayList<int>



Valhalla:

- Inline types
- Specialized generics
- Migration

Identity, Value objects and other theories.

Identity

Identity (object-oriented programming)

From Wikipedia, the free encyclopedia

An **identity** in object-oriented programming, object-oriented design and object-oriented analysis describes the property of objects that distinguishes them from other objects. This is closely related to the philosophical concept of <u>identity</u>.

In philosophy, **identity**, from Latin: *identitas* ("sameness"), is the relation each thing bears only to itself. The notion of identity gives rise to many philosophical problems, including the identity of indiscernibles, and questions about change and personal identity over time.

Identity

The idea is the following: in a model with object identity, an object has an existence which is independent of its value. Thus two notions of object equivalence exist: two objects can be identical (they are the same object) or they can be equal (they have the same value). This has two implications: one is object sharing and the other one is object updates.

Value object

Value object

From Wikipedia, the free encyclopedia

In computer science, a **value object** is a small object that represents a *simple* entity whose <u>equality</u> is not based on identity: i.e. two value objects are *equal* when they *have* the same *value*, not necessarily being the *same object*.^{[1][2]}

ValueObject

14 November 2016



Martin Fowler

https://martinfowler.com/bliki/ValueObject.html

Identity and mutability

Identity enables mutability^(*)

(*) in order to mutate a field of an object, we must know which object we are trying to modify

Value object and mutability

http://wiki.c2.com/?ValueObject

http://wiki.c2.com/?ValueObjectHypotheses

http://wiki.c2.com/?ValueObjectsCanBeMutable

http://wiki.c2.com/?ValueObjectsShouldBeImmutable

http://wiki.c2.com/?CanValueObjectsContainReferenceObjects

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Java

C#

Identity

The "simplest" identity implementation – reference/pointer/address

Identity gives:

• Extra value null

Extra attributes

"each object in Java is associated with a monitor" (c) JLS

Performance

e.g. $(x==y \mid | x.equals(y))$

Identity gives:

Not always working optimizations

e.g. Escape analysis + scalar replacement

Not optimal optimizations

e.g. -XX:+UseStringDeduplication

Extra overhead to preserve identity

e.g. ZGC/Shenandoah read barriers

Does Java already has Value Objects?



Does Java already has Value Objects(*)?

- java.lang.String
- java.lang.Integer

•

(*) - theoretically we may deprive these classes of identity

Does Java already has Value Objects?

Value-based Classes

Some classes, such as java.util.Optional and java.time.LocalDateTime, are *value-based*. Instances of a value-based class:

- are final and immutable (though may contain references to mutable objects);
- have implementations of equals, hashCode, and toString which are computed solely from the instance's state and not from its identity or the state of any other object or variable;
- make no use of identity-sensitive operations such as reference equality (==) between instances, identity hash code of instances, or synchronization on an instances's intrinsic lock;
- are considered equal solely based on equals(), not based on reference equality (==);
- do not have accessible constructors, but are instead instantiated through factory methods which make no committment as to the identity of returned instances;
- are *freely substitutable* when equal, meaning that interchanging any two instances x and y that are equal according to equals() in any computation or method invocation should produce no visible change in behavior.

Does Java already has Value Objects?

This implies that the identity of the result of evaluating a lambda expression (or, of serializing and deserializing a lambda expression) is unpredictable, and therefore identity-sensitive operations (such as reference equality (§15.21.3), object locking (§14.19), and the System.identityHashCode method) may produce different results in different implementations of the Java programming language, or even upon different lambda expression evaluations in the same implementation.

What is Inline class?



Inline Class

```
public inline class Complex {
    private double re, im;
    public Complex(double re, double im) {
        this.re = re;
        this.im = im;
    public double re() { ... }
    public double im() { ... }
    public Complex add(Complex c) { ... }
    public Complex mult(Complex c) { ... }
```

Inline Class

```
public(inline) class Complex {
    private double re, im;
    public Complex(double re, double im) {
        this.re = re;
        this.im = im;
    public double re() { ... }
    public double im() { ... }
    public Complex add(Complex c) { ... }
    public Complex mult(Complex c) { ... }
```

No Identity

Inline class

- is a class
- no identity → immutable, non-nullable
- that's all

We didn't say:

- it's allocated on heap or out of heap
- it's allocated on stack or out of stack
- it's inlined into container object or referenced by container object

Identity-sensitive operations

Synchronization (Object::wait, Object::notify)

throw new IllegalMonitorStateException();

Identity-sensitive operations

Equality (==), substitutability check

```
a == b
```

- a & b has the same type
- for each field:
 - -- if inline class check substitutability
 - -- if double/float check Double/Float::equals
 - -- else check "=="

Identity-sensitive operations

- System::identityHashCode
 - compute over fields

Identity-sensitive operations

Weak references

JDK-8244968 [valhalla] WeakHashMap needs to handle keys which are inline objects

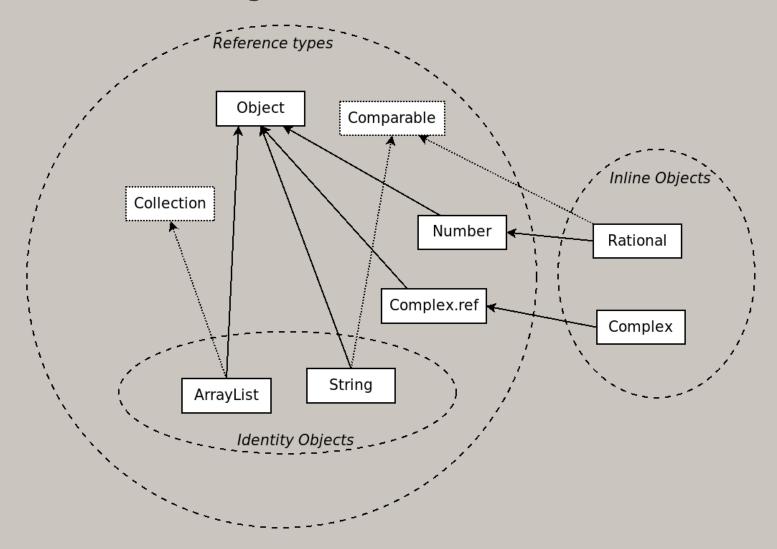
Class hierarchy



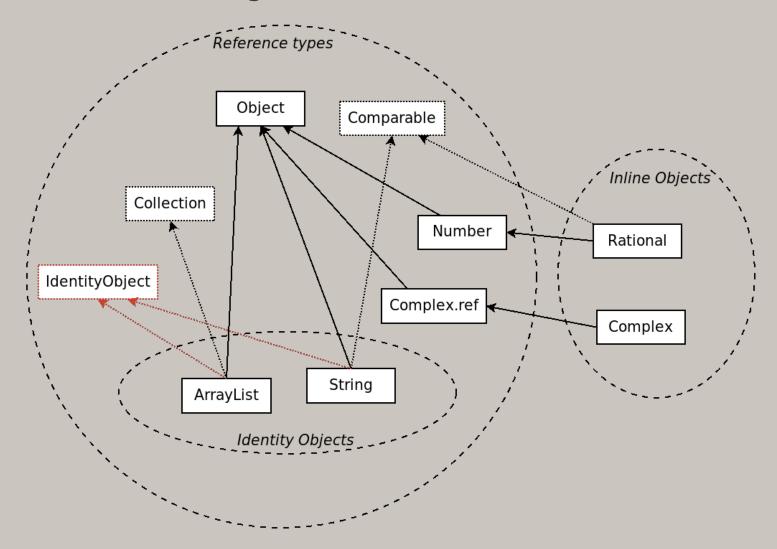
Inline class

- is final
- can implement interfaces
- extends "well-formed" abstract class:
 - no fields
 - empty no-arg constructor
 - no synchronized methods
 - has "well-formed" super

Class hierarchy



Class hierarchy



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IdentityObject

```
if (x instanceof IdentityObject) {
    synchronized(x) { ... }
}

void runWithLock(IdentityObject lock, Runnable r) {
    synchronized (lock) {
        r.run();
    }
}
```

0

Wait, wait!

What about "new Object()"?

Wait, wait!

What about "new Object()"?

@Deprecated

Wait, wait!

What about "new Object()"?

@Deprecated

```
class java.util.Objects {
    /**
    * Constructs a new Object implementing the IdentityObject interface.
    * The object is a unique IdentityObject suitable for all purposes
    * that previously for which new Object() was used including
    * synchronization, mutexes and unique placeholders.
    *
     * @return a new Object implementing the IdentityObject interface
     * @since Valhalla
     */
    public IdentityObject newIdentity()
     ...
```

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What about arrays?

The same as before:

W

Boxing
vs
Inline Widening Conversion



Value Set (pre Valhalla)

primitive type → primitive values

reference type → null + references to object instances

Value Set (Valhalla)

```
primitive type → primitive values
```

```
inline type → object instances
```

identity type → null + references to object instances

reference type → null + ???

Value Set (Valhalla)

```
primitive type → primitive values
```

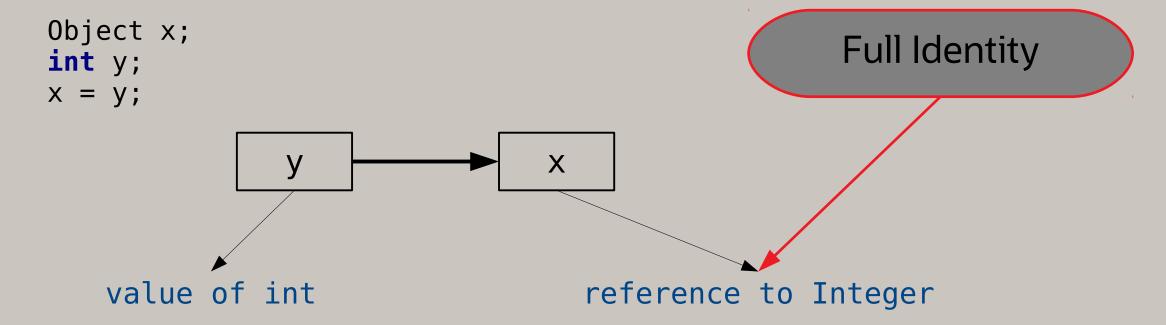
```
inline type → object instances
```

identity type → null + references to object instances

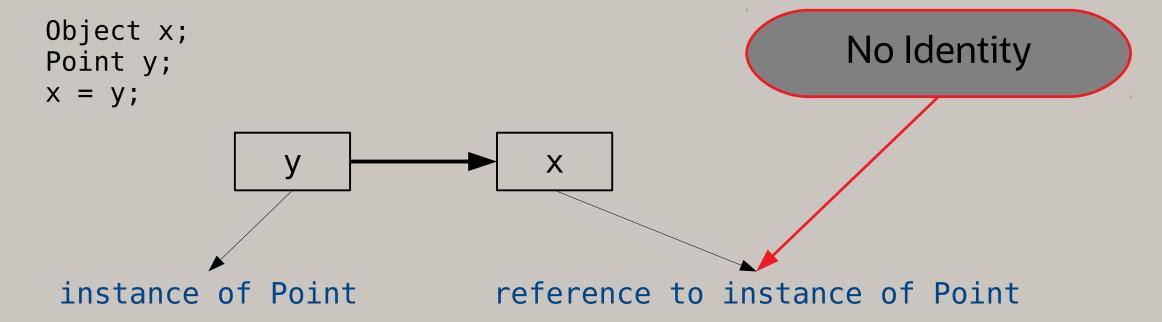
reference type → null + references to

both identity and inline objects

Boxing



Inline widening



No more questions



T – inline type

T. ref – reference projection for T

ValSet(T.ref) = { null } \cup { ref v : v \in ValSet(T) }

inline class V {...}

sealed abstract class V.ref permits V.val { }

We generated
inline class V.val extends V.ref {...}

T – inline type

T. ref – reference projection for T

T. val – value projection for T

T.ref → T.val – *inline narrowing conversion*

S – reference type

$$S.ref == S$$

Default value



Default value

T – inline type

T.default - default value of T

for each field:

-- 0, 0.0, false, default or null

Default value

S – reference type

S.default == null

Implementation details



If you (your JVM) are very lazy

- Allocate all inline classes in heap (as other objects)
- Implement all identity-sensitive operations in a proper way
- Don't forget to make inline types non-nullable
- Some boring stuff with classloading and verification

What about Hotspot?

- Sometimes allocate inline classes in heap
- Inline it into containers (arrays, classes)
- Locals and parameters scalarized

Migration stories



One migration story

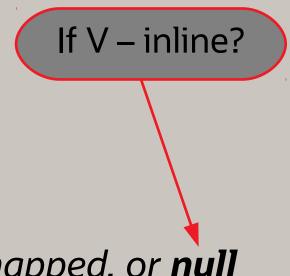
```
Map<K,V>
V get(Object key);
```

Returns the value to which the specified key is mapped, or **null** if this map contains no mapping for the key.

One migration story

```
Map<K,V>
V get(Object key);
```

Returns the value to which the specified key is mapped, or **null** if this map contains no mapping for the key.



One migration story

```
Map<K,V>
V.ref get(Object key);
```

Returns the value to which the specified key is mapped, or **null** if this map contains no mapping for the key.

Optional<T>

API Note:

Optional is primarily intended for use as a method return type where there is a clear need to represent "no result," and where using null is likely to cause errors. A variable whose type is Optional should never itself be null; it should always point to an Optional instance.

```
Optional<T> - will be inline
but:
Optional<T> op;
...
op = null;
Can't be inline
```

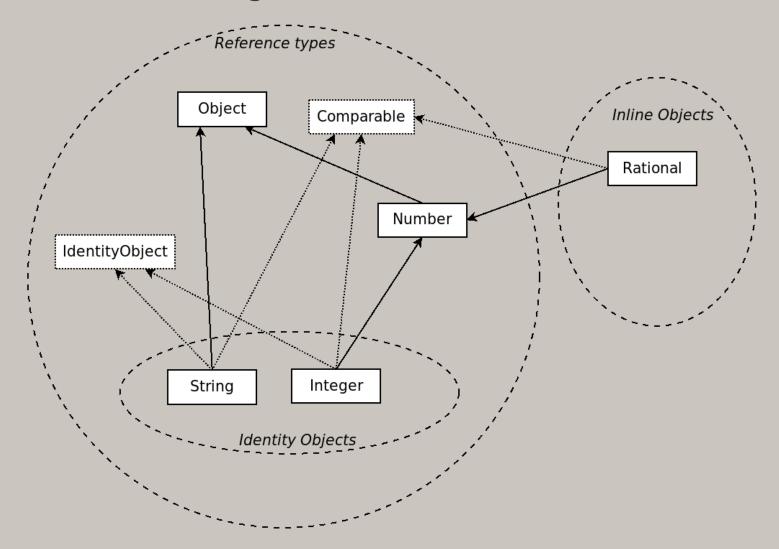
```
inline class V {...}
V x; ... foo(V a); ... T<V> ...; // usages of V
                                 // V - just an alias
sealed abstract class V.ref permits V.val { }
inline class V.val extends V.ref {...}
```

inline class V.val extends V.ref {...}

Very primitive migration story



Class hierarchy



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We forgot primitives

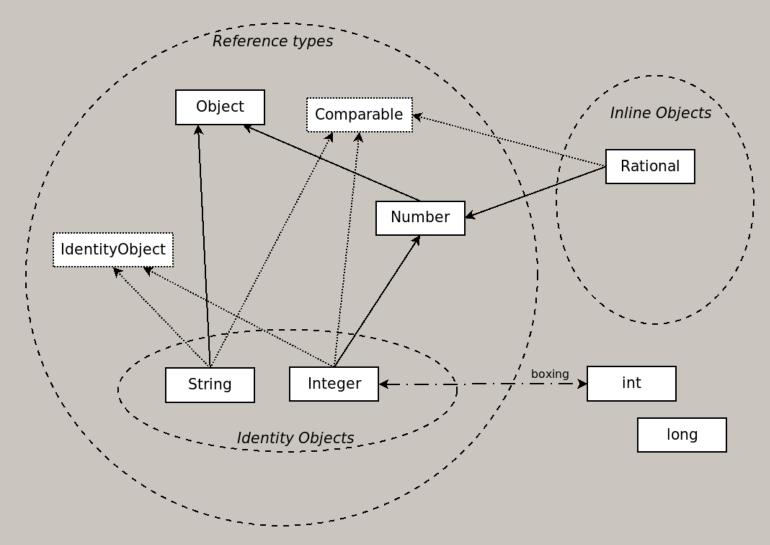
- Not a true OOP
- Incompatible arrays, no generics
- Complex code:

```
e.g. Arrays::fill - 9 methods
```

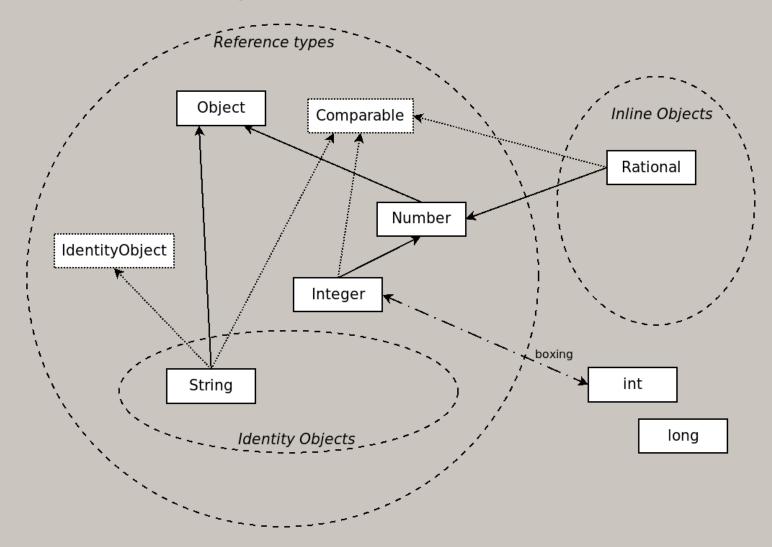
Complex libraries:

e.g. Stream<T>, IntStream, LongStream, DoubleStream

Type hierarchy

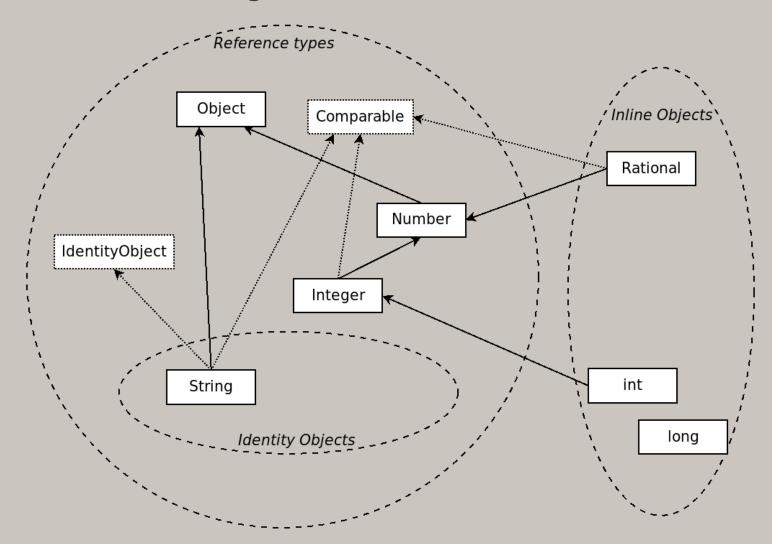


Type hierarchy



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Class hierarchy



First step

Deprecate identity of primitive wrappers

JDK-8236567 diagnostic anti-locked objects, dynamically by instance

JDK-8236568 diagnostic anti-locked objects, by class

JDK-8242263 Diagnose synchronization on primitive wrappers

Second step

• Implement:

```
inline class int { ... }
```

- Integer == int.ref
- Integer.val == int

As result

- Make Java true OOP again
- Fully covariant arrays:

```
int[] <: Integer[] <: Object[]</pre>
```

Simplify generic specialization:

As result

at the beginning:

Inlines – user-defined primitives

at the end:

Primitives – built-in inlines

Collateral benefits



Nestmates

JEP 181: Nest-Based Access Control

@since 11

https://openjdk.java.net/jeps/181

Field layout computation overhaul

JDK-8237767 Field layout computation overhaul

@since 15

-XX:+UseEmptySlotsInSupers

Real objects size (in bytes)

		+UseCompressedOops		-UseCompressedOops	
		Java 14	Java 15	Java 14	Java 15
class L1B0	{ long l1; }	24	24	24	24
class L1B1 extends L1B0) { byte b1; }	32	24	32	32
class L2B1 extends L1B1	l { long l2; }	40	32	40	40
class L2B2 extends L2B1	l { byte b2; }	48	32	48	40
class L3B2 extends L2B2	2 { long l3; }	56	40	56	48
class L3B3 extends L3B2	2 { byte b3; }	64	40	64	48
class L4B3 extends L3B3	3 { long l4; }	72	48	72	56
class L4B4 extends L4B3	3 { byte b4; }	80	48	80	56

Specialized generics



Erased Generics

"Background: how we got the generics we have

(Or, how I learned to stop worrying and love erasure)"

Brian Goetz, June 2020

http://cr.openjdk.java.net/~briangoetz/valhalla/erasure.html

Erased Generics

"Two Ways to Bake Your Pizza —

Translating Parameterised Types into Java"

Martin Odersky, Enno Runne, and Philip Wadler, 2000

http://pizzacompiler.sourceforge.net/doc/pizza-translation.pdf

Specialization

- Layout specialization
 - •List<int> no boxing
- Polymorphic specialization
 - Predicate<T> == Function<T, boolean>
 - HashSet<T> == HashMap<T, void>

Reification

- Reified generics preserve full type information
 - (x instance of List<T>)
 - (List<int>)x;

Generics

- Java has erased generics → Java will preserve it
- Focus on specialization
 - List<reference type> erased
 - List<inline type> specialized
- Reification not a goal

Plans

- Two phases:
 - expand expressiveness (List<int>)
 - layout specialization
- No compile time specialization
 - only at runtime and only by JVM decision
- To be continued....





- "Complex[size]" vs "double[2*size]"
- Which layout? Native interop?

- "Complex[size]" vs "double[2*size]"
- Which layout? Native interop?

- Valhalla doesn't specify layout
- @see Panama

Fused String, fused arrays (inject array into object)?

Fused String, fused arrays (inject array into object)?

- Valhalla doesn't do fused arrays
- @see Arrays 2.0

Why did you take name which has different meaning in Kotlin?

Why did you take name which has different meaning in Kotlin?

Because we can.

Links

"State of Valhalla";

- 1. The Road to Valhalla
- 2. Language Model
- 3. JVM Model
- 4. Translation scheme

https://cr.openjdk.java.net/~briangoetz/valhalla/sov/01-background.html

Links

Wiki:

https://wiki.openjdk.java.net/display/valhalla/Main

Mailing lists:

http://mail.openjdk.java.net/mailman/listinfo/valhalla-dev

http://mail.openjdk.java.net/mailman/listinfo/valhalla-spec-observers

Repository:

https://github.com/openjdk/valhalla

Thank You

Sergey Kuksenko

Java Platform Group Oracle

