Propositional Logic

Aiphabet - aiphabet.org

Logic is all about how to conceptualize and concisely write concepts and statements, and make inference/draw conclusions from them. Logic comes in different flavors. The simplest one being Propositional Logic (PL). The building blocks of PL are propositions and connectives.

Definition 1. *Proposition:* A proposition is a declarative statement that's either true or false. For example, "it is hot today", "2+3=4", "if it is hot then we go swim" are propositions.

In general, statements that are questions, commands, or opinions are not propositions. For example "Hello!", "how are you", "give me the key", are not propositions.

The **building blocks** of PL are propositions and connectives. The simplest form of a proposition is called an **atomic proposition**. We use lower-case letters to denote atomic propositions, such as p, q, r, t, s, etc.

In natural language, we make statements and combine statements with words, such as "and," "or," "if-then," "but," etc. In natural language, we make statements and combine them with words, such as "and", "or," "if-then," "but," etc. Likewise, in PL, we create compound propositions by combining atomic propositions using logical connectives.

1 Truth Tables

Before we review the major connectives, we define a useful table that allows us to evaluate the truth of a compound proposition based on the values of its atomic propositions.

• Negation: Let p be a proposition. We call "not p" the negation of p, denoted by $\neg p$.

p	$\neg p$
т	F
F	т

Conjunction: Let p and q be two propositions. We call conjunction the compound proposition p and q, denoted by p ∧ q.

p	q	
т	т	т
т	F	F
F	т	F
F	F	F

• **Disjunction:** Let p and q be two propositions. We call disjunction the compound proposition p or q, denoted by $p \lor q$.

p	q	$p \lor q$
т	т	т
т	F	т
F	т	т
F	F	F

• Implication (if-then): We can implication the compound proposition "if p then q", denoted $p \rightarrow q$ or $p \Rightarrow q$.

p	q	$p \rightarrow q$
т	т	т
т	F	F
F	т	т
F	F	т

If and only if (IFF): Let p and q be two propositions. We call an equivalence the compound proposition "p if and only if q", denoted as p ↔ q or p ⇔ q.

p	q	$p \leftrightarrow q$
т	т	т
т	F	F
F	т	F
F	F	т

Just like arithmetic operators, there is a precedence order when evaluating logical connectives as follows (from highest to lowest): () processed from inside to outside, $\neg, \wedge, \lor, \rightarrow, \leftrightarrow$, left to right.

Definition 2. *Tautology*: A tautology is a proposition that is always true for any assignment of truth values to its atomic propositions.

Definition 3. *Fallacy/Contradiction*: A fallacy (or a contradiction) is a proposition that is always false for any assignment of truth values to its atomic propositions.

Definition 4. *Logical Equivalence*: *Two propositions p and q are logically equivalent if and only if they have the same truth tables.*

2 Laws of Propositional Logic

1. Identity law:

$$p \wedge True \equiv p$$

 $p \vee False \equiv p$

- 2. Domination law:
 - $p \lor True \equiv True$

 $p \wedge False \equiv False$

3. Idempotence law:

$$p \vee p \equiv p$$

 $p \wedge p \equiv p$

4. Negation law:

$$p \land (\neg p) \equiv False$$

$$p \lor (\neg p) \equiv True$$

5. Double negation law:

$$\neg \neg p \equiv p$$

6. Commutativity law:

$$p \wedge q \equiv q \wedge p$$

 $p \lor q \equiv q \lor p$

3 Logical Inference

Definition 5. *Inference/Argument*: An inference in Propositional Logic is a sequence of propositions denoted as:

$$\frac{p_1 \ p_2 \dots p_n}{q}$$

where p_1, p_2, \ldots, p_n are called **premises** and q is called **conclusion**.

Definition 6. The inference is called valid if premises p_1, \ldots, p_n implies q, in other words

 $(p_1 \wedge p_2 \wedge \ldots \wedge p_n) \rightarrow q$ is a tautology

MInference and implication:

An inference rule is an implication that is always true (tautology).

7. Associativity law:

 $(p \land q) \land r \equiv p \land (q \land r)$ $(p \lor q) \lor r \equiv p \lor (q \lor r)$

- 8. Distributivity law:
 - $p \land (q \lor r) \equiv (p \land q) \lor (p \land r)$ $p \lor (q \land r) \equiv (p \lor q) \land (p \lor r)$ $p \lor (q \lor r) \equiv (p \lor q) \lor (p \lor r)$ $p \land (q \land r) \equiv (p \land q) \land (p \land r)$
- 9. Absorption law:

$$p \lor (p \land q) \equiv p$$

$$p \land (p \lor q) \equiv p$$

10. DeMorgan's law:

$$\neg (p \land q) \equiv (\neg p) \lor (\neg q)$$
$$\neg (p \lor q) \equiv (\neg p) \land (\neg q)$$

11. Implication to disjunction law:

 $p \to q \equiv \neg p \lor q$

12. IFF to implication law:

 $p \leftrightarrow q \equiv (p \rightarrow q) \land (q \rightarrow p)$

Example 1:

$$\frac{\neg \neg p}{p}$$
 is a valid inference.

We can prove this using a truth table.

$\neg \neg p$	р	$\neg \neg p ightarrow p$
т	Т	т
F	F	т

3.1 Rules of Inference

V I	Rules of inference			
	Addition	Conjunction	Simplification	Disjunctive syllogism
	p	p q	$p \wedge q$	$p \lor q \neg p$
	$\overline{p \lor q}$	$\overline{p \wedge q}$	\overline{p}	\overline{q}
	Hypothetical syllogism	Resolution	Modus ponens	Modus tollens
	$p \rightarrow q q \rightarrow r$	$\underline{p \lor q \neg p \lor r}$	$\underline{p p \rightarrow q}$	$\underline{\neg q p \to q}$
	$p \rightarrow r$	$q \lor r$	q	$\neg p$

Example 2: Use of modus ponens:

 If you forget your jacket then you will be cold
 You forgot your jacket today

 You will be cold
 You will be cold

$$\frac{p \to q}{q} p$$